# MAT 145: Homework 6 Solution 

Xiaochen Liu*

1. One possible answer is $A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F \rightarrow G \rightarrow H \rightarrow A$.

2. One possible example is $K_{4}$.
3. One possible example is $K_{3}$.
4. Existence is ensured by connectness of $G$. Now we prove uniqueness. Assume that $v_{i}$ and $v_{j}$ are connected by at least two different paths. For example, $v_{i} \rightarrow a_{1} \rightarrow \cdots \rightarrow a_{n} \rightarrow v_{j}$ and $v_{i} \rightarrow b_{1} \rightarrow \cdots \rightarrow b_{m} \rightarrow v_{j}$ are two different paths. Along the path of $a$ 's from $v_{i}$ to $v_{j}$, find the first vertex where it diverges from the path of $b$ 's, i.e., the smallest $k$ and $\ell$ such that $a_{k}=b_{\ell}$ and $a_{k+1} \neq b_{\ell+1}$. We continue walk along the path of $a$ 's, find the first vertex where the two paths meet again, i.e., the smallest $p>k$ such that $a_{p-1} \neq b_{q-1}$ and $a_{p}=b_{q}$. Then there are two disjoint paths connecting $a_{k}\left(=b_{\ell}\right)$ and $a_{p}\left(b_{q}\right)$. So, there is a cycle in $G$. A contradiction.
5. Suppose $G^{\prime}$ is obtained by adding an edge $u v$ to $G$ and $u v \notin E(G)$. In $G, u$ and $v$ are already connected by a path as $G$ is connected. So, there are two disjoint paths connecting $u$ and $v$, thus there is a cycle.
6. Consider the six 4 -vertex trees on page 145 . For the leftmost one, for example, call it $G_{1}$. Add a new vertex to $G_{1}$. This new vertex may connect to any of the four existing vertices. So $G_{1}$ has four 'offspring'. Do the same thing for the other five trees. Figures omitted.

[^0]7. Let $e$ be the number of edges in $T, v$ be the number of vertices in $T, v_{j}$ be the number of vertices with degree $j$ in $T$, for $j=1,3$. Then we have $e=v-1, e=2 v=v_{1}+3 v_{3}, v_{1}+v_{3}=v$ and $v_{3}=10$. Solving this linear system gives $v_{1}=12$.
8. One possibility is as follow.

9. 0103553.
10. 310442552.


[^0]:    *xchliu@math.ucdavis.edu

