## Final Exam

Math 145, Spring 2019

## Name: Solutions

## Student ID:

Every solution must contain an explanation written in words supporting your numerical solution to reccive credit.

You do not need to simplify numerical expressions for your final answers (e.g. you can write $2^{8} \cdot 4$ ! instead of multiplying out to 6144.)

If you need extra space for your solutions, there is an extra page at the back of the exam. If you need extra space for any problem, write CONTINUED IN EXTRA SPACE on the page where the problem is given to you. In the extra space write the problem number that you are solving in that space.

Problem 1: Let $T$ be a tree. Prove that every edge of $T$ is a cut-edge.

Let $e$ be any edge of $T$. Let $v_{i}$ and $v_{j}$ be the endpoints of $e$.
Let $T^{\prime}$ be the graph obtained from $T$ by deleting $e$, from the edge set. Suppose for contradiction that $T^{\prime}$ was connected.
Then the ne exists a linear subgraph of $T^{\prime}$ connecting $v_{i}$ to $v_{j}$


Sine $e$ is not on edge of $T^{\prime}$, $e$ is not included in this linear subgraph.
Then the union of the linear subgraph with the edge $e$ is a subgraph of $T_{1}$, but this subgraph is a cycle

but $T$ is a tree so it cannot contain a cycle.

Problem 2: How many numbers between 1 and 100 are NOT a multiple of any of the numbers $3,5,11$ ?
Define $A, B, C \subseteq\{1,2, \ldots, 100\}$
(arithmetic migtaness were forgive)
ky:

$$
\begin{array}{ll}
A=\{\text { multiples of } 3\} & \left.|A|=L^{100} / 3\right\rfloor=33 \\
B=\{\text { multiples of } 5\} & |B|=\left\lfloor^{100} / 5\right\rfloor=20 \\
C=\{\text { multiples of } 11\} & \left.|C|=L^{100} / 11\right\rfloor=9
\end{array}
$$

$A \cap B=\{$ multiples of $\left.3.5 \cdot 15\} \quad|A \cap B|=L^{100} / 15\right\rfloor=6$
$A \cap C=\{$ multiples of $\left.3 \cdot 11=33\} \quad|A \cap C|=L^{100 / 33}\right\rfloor=3$
$B \cap C=\{$ multiples of $511=55\} \quad|B \cap C|=\lfloor 100 / 55\rfloor=1$
$A \cap B \cap C=\{$ multiples of $3.5 \cdot 11=165\}=\varnothing \quad|A \cap B \cap C|=0$
Then numbers which are a multiple of some number $3,5,11$ are

$$
\begin{aligned}
|A \cup B \cup C| & =|A|+|B|+|C|-|A \cap B|-|A \cap C|-|B \cap C|+|A \cap B \cap C| \\
& =33+20+9-6-3-1+0 \\
& =52
\end{aligned}
$$

So the numbers which are not a multiple of amy of muse one $\{1,2, \ldots, 100\} \backslash(A \cup B \cup C)$
so there are $100-52=48$ such numbers.

Problem 3: For the following pair of graphs, either prove they are isomorphic or prove they are not isomorphic.

Gi


We construct an isomorphism


$$
\begin{array}{ll}
\Phi\left(v_{1}\right)=w_{1} & \psi\left(e_{1}\right)=f_{2} \\
\Phi\left(v_{2}\right)=w_{4} & \psi\left(e_{2}\right)=f_{1} \\
\Phi\left(v_{5}\right)=w_{2} & \psi\left(e_{3}\right)=f_{4} \\
\Phi\left(v_{4}\right)=w_{5} & \psi\left(e_{4}\right)=f_{3} \\
\Phi\left(v_{5}\right)=w_{3} & \psi\left(e_{5}\right)=f_{5}
\end{array}
$$

$G_{2}$


| endpoints in $G_{1}$ | endpoints in $G_{2}$ |
| :--- | :--- |
| $e_{1}: v_{1}, v_{3}$ | $\Psi\left(e_{1}\right)=f_{2}: w_{1}, w_{2}=\Phi\left(v_{1}\right), \Phi\left(v_{3}\right)$ |
| $e_{2}: v_{1}, v_{4}$ | $\psi\left(e_{2}\right)=f_{1}: w_{1}, w_{5}=\Phi\left(v_{1}\right), \Phi\left(v_{4}\right)$ |
| $e_{3}: v_{2}, v_{5}$ | $\psi\left(e_{3}\right)=f_{4}: w_{3}, w_{4}=\Phi\left(v_{5}\right), \Phi\left(v_{2}\right)$ |
| $e_{4}: v_{3}, v_{5}$ | $\psi\left(e_{4}\right)=f_{3}: w_{2}, w_{3}=\Phi\left(v_{3}\right), \Phi\left(v_{5}\right)$ |
| $e_{5}: v_{4}, v_{2}$ | $\psi\left(e_{5}\right)=f_{5}: w_{4}, w_{5}=\Phi\left(v_{2}\right), \Phi\left(v_{4}\right)$ |

other solutions:

Problem 4: Let $T$ be the labeled tree that corresponds to the Prüfer code 47274072. (a) Draw the tree $T$ and (b) write down the degree sequence for the vertices of $T$.
(a)

$$
\frac{1}{4} \frac{3}{7} \frac{5}{2} \frac{6}{7} \frac{8}{4} 4 \frac{9}{0} 72 \frac{2}{0}
$$


(b) Degrees:

| Vertex | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Degree | 2 | 1 | 3 | 1 | 3 | 1 | 1 | 4 | 1 | 1 |

Ordered: $(4,3,3,2,1,1,1,1,1,1)$

Problem 5: Suppose $G$ is a simple connected graph with a planar embedding that cuts up the plane into 8 regions. If $G$ has 12 vertices,
(a) Calculate (with proof) the sum of the degrees of all the vertices.

$$
\sum_{v \in V(G)} d(v) .
$$

Euler's formula: \#V-\#E+\#R=2

$$
\begin{gathered}
12-H E+8=2 \\
\Rightarrow \quad 18=\# E
\end{gathered}
$$

Edge-degree formula: $\sum_{v \in V(G)} d(v)=2 \neq E=36$
(b) Next, draw an example of a planar embedding of a graph $G$ with these properties ( 8 regions in the planar embedding and 12 vertices).


Problem 6: Give the proof for the statement that if a graph $G$ has a closed Eulerian walk, then every vertex of $G$ has even degree.

An Eulerion walk uses every edge exactly on a.
If it is a closed Eulecian walk, the starting + ending vertex are the same.
Therefore every vertex is entered on the walk the same number of times it is exited.

Since every edge must be used, every endpoint of ever edge must be used either as on entering point or on exiting point.

The degree of a vertex $v$ is the total number of endpoints of edges that agree with $v$.
Since half of these endpoints must be entering ports and the other half must- be exiting points the total number of endpoints of edges at $r$ must be even. Therefor the degree at $v$ must be ever.

Problem 7: Carla is choosing the dinner buffet for a celebration. First, she must choose whether to serve 3,4 or 5 different main course dishes. Once she decides how many dishes to serve, she must pick the dishes from a list of 12 main course options. Then she has to choose 2 dessert choices from a list of 6 dessert options. How many different possible dinner and dessert combination selections does Carla have? (We consider all dinners with 3 main course options different from all dinners with 4 main course options, etc.)


Given $n$ options, we wort to choose $k$ of them, unordered, non repeating so the number of such options is $\binom{n}{k}$.

Problem 8: Find the minimal cost spanning tree for the following weighted graph. Draw the spanning tree and determine its total cost.

$\left.\begin{array}{l}1 \text { Include } \\ 3 \\ 3 \\ \text { Include } \\ 3 \\ \text { Include } \\ \text { Include }\end{array}\right\} \rightarrow \begin{array}{lll}9 & 3 \\ 3 & & \text { no cycles }\end{array}$ 4 exclude because could create a cycle 5 exclude bl a would create a cycle 6 Include

7 exclude-cycle

$\left.\begin{array}{l}8 \text { Include } \\ 8 \text { Include } \\ 9 \text { Include }\end{array}\right\rangle$


Now all vertices ane connected so we stop adding edges
Total cost: $1+2+3+3+6+8+8+9=40$

Problem 9: Suppose $G$ is a simple graph with 8 vertices and 17 edges. Prove that $G$ has at least one vertex of degree strictly greater than 4 .

Degree-edge formula:

$$
\sum_{\substack{v \in V(G) \\ ק}} d(v)=2 \# E(G)=2.17=34
$$

8 vertices

$$
d\left(v_{1}\right)+d\left(v_{2}\right)+\ldots+d\left(v_{p}\right)=34
$$

If every degree was less than or equal to 4 then

$$
d\left(v_{1}\right)+d\left(v_{2}\right)+\ldots+d\left(v_{8}\right) \leq \underbrace{4+4+\ldots+4}_{8 \text { times }}=4.8=32
$$

But 34 is not $\leq 32$
So we get a contradiction.
Thus there must be at least one vertex with degree greater then 4 .

Extra Space:

