

Low Rank Approximations

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Thm For any k with $0 \leq k \leq r$,
let $A_k := \sum_{j=1}^k \sigma_j u_j v_j^T$

If $k = p = \min(m, n)$, then define $\sigma_{k+1} = 0$. Then,

$$\|A - A_k\|_2 = \inf_{\substack{B \in \mathbb{R}^{m \times n} \\ \text{rank}(B) \leq k}} \|A - B\|_2 = \sigma_{k+1}$$

(Proof)

$$\|A - A_k\|_2 = \left\| \sum_{j=k+1}^p \sigma_j u_j v_j^T \right\|_2$$

$$= \left\| U \begin{bmatrix} \sigma_{k+1} & & 0 \\ & \ddots & \\ 0 & & \sigma_p \\ \hline & & 0 \end{bmatrix} V^T \right\|_2$$

$$= \left\| \begin{bmatrix} \sigma_{k+1} & & 0 \\ & \ddots & \\ 0 & & \sigma_p \\ \hline & & 0 \end{bmatrix} \right\|_2$$

since
 U, V :
orthogonal!

$$= \sigma_{k+1} \quad \text{by definition of the matrix norm}$$

Note: If $D = \text{diag}(d_1, \dots, d_m) = \begin{bmatrix} d_1 & & 0 \\ & \ddots & \\ 0 & & d_m \end{bmatrix}$
then $\|D\|_p = \max_{1 \leq j \leq m} |d_j| \quad \forall p \geq 1$

Prove
this $\Rightarrow$
as an
exercise!

Now, let $B \in \mathbb{R}^{m \times n}$ be any rank k matrix. Then $\dim(\text{null}(B)) = n - k$
why? Because of the following Thm:

For any $A \in \mathbb{R}^{m \times n}$, $\text{rank}(A) + \dim(\text{null}(A)) = n$

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Let $W := \text{null}(B) \cap \langle v_1, \dots, v_{k+1} \rangle$

We know $W \neq \{0\}$ because

$$\dim(\text{null}(B)) = n - k$$

$$\dim(\langle v_1, \dots, v_{k+1} \rangle) = k+1$$

So, if these two do not intersect, $\mathbb{R}^n$'s dimension would become $n - k + k + 1 = n + 1$

This cannot happen! #

So let $th \in W$, $th \neq 0$.

We can always normalize th , so can assume $\|th\|_2 = 1$.

Then,

$$\|A - B\|_2 \geq \|(A - B)th\|_2 \text{ by def.}$$

$$= \|Ath\|_2 \text{ since } th \in \text{null}(B)$$

$$= \|U \Sigma V^T th\|_2$$

Since $th \in \langle v_1, \dots, v_{k+1} \rangle$

$$V^T th = \begin{bmatrix} * \\ \vdots \\ * \\ 0 \\ \vdots \\ 0 \end{bmatrix} \begin{matrix} \left. \vphantom{\begin{bmatrix} * \\ \vdots \\ * \\ 0 \\ \vdots \\ 0 \end{bmatrix}} \right\} k+1 \\ \left. \vphantom{\begin{bmatrix} * \\ \vdots \\ * \\ 0 \\ \vdots \\ 0 \end{bmatrix}} \right\} n-k-1 \end{matrix}$$

$$= \|\Sigma V^T th\|_2 \text{ since } U: \text{ortho.}$$

$$\geq \sigma_{k+1} \|V^T th\|_2$$

$$= \sigma_{k+1} \|th\|_2 = \sigma_{k+1}$$

///

Thm For any k with $0 \leq k \leq r$,

$$\|A - A_k\|_F = \inf_{\substack{B \in \mathbb{R}^{m \times n} \\ \text{rank}(B) \leq k}} \|A - B\|_F$$

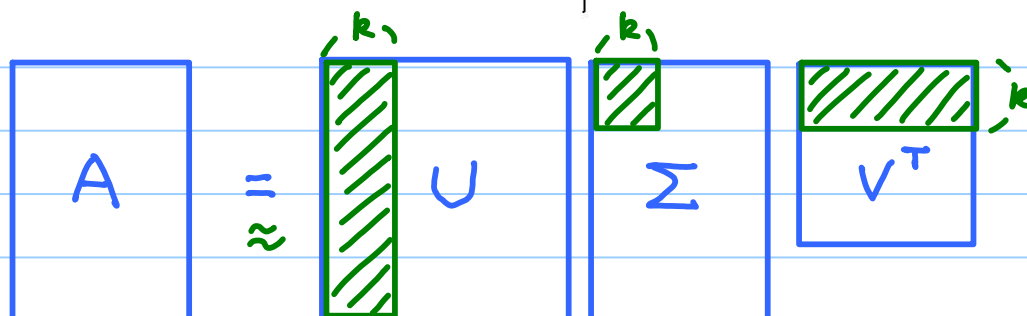
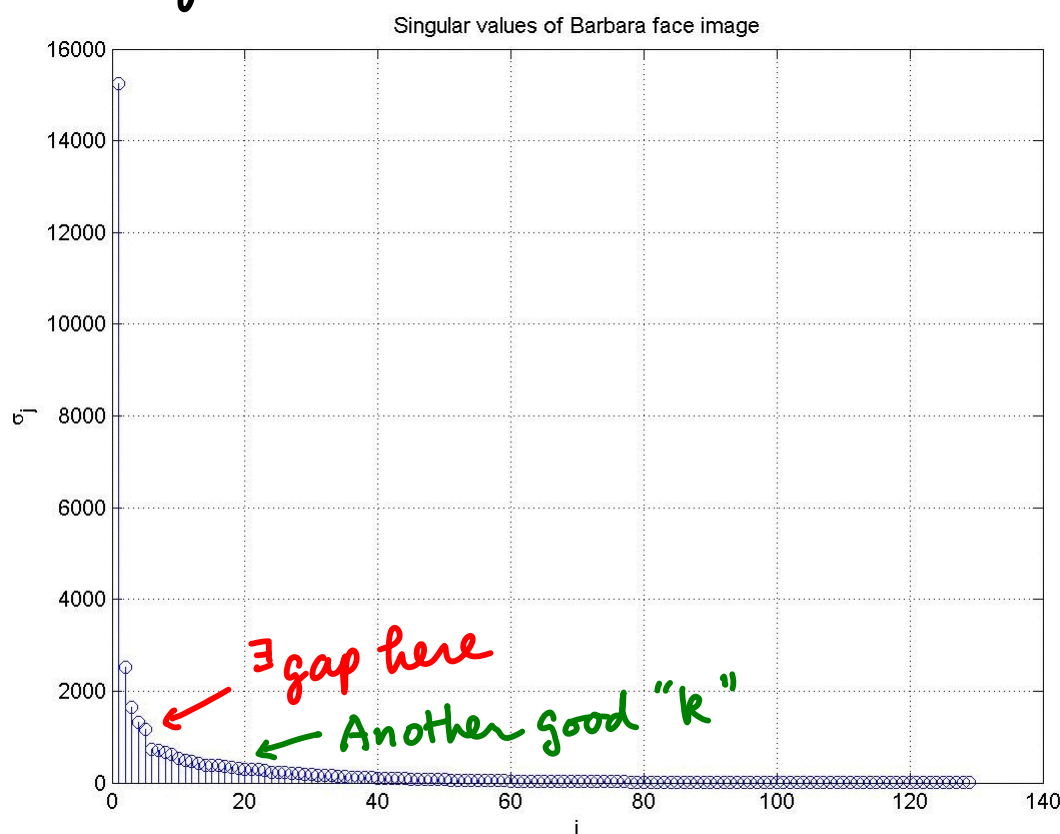
$$= \sqrt{\sigma_{k+1}^2 + \dots + \sigma_r^2}$$

(Proof) Exercise!

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So, for a given matrix, say, A
how to determine a good "k"
so that we can efficiently (i.e.,
compress) A without losing too much
info of A ?

⇒ Check the distribution of the
singular values!



rank k approximation of A only uses $|||||$ portions!

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★ Condition Number and SVD

Recall the condition number for a square nonsingular matrix A :

$$\kappa(A) = \text{cond}(A) := \|A\|_2 \|A^{-1}\|_2$$

$\kappa(A)$: small $\Rightarrow A$: well-conditioned.

$\kappa(A)$: large $\Rightarrow A$: ill-conditioned,
lose $\approx \log_{10} \kappa(A)$ digits
to solve $Ax = b$.

If A : singular, $\kappa(A) = +\infty$.

Using SVD of A , we can nicely compute $\kappa(A)$ as follows.

$$\|A\|_2 = \sigma_1 \quad \rightarrow \text{by definition}$$

$$\|A^{-1}\|_2 = 1/\sigma_m \quad \text{why?} \quad A^{-1} = (U \Sigma V^T)^{-1} = V \Sigma^{-1} U^T \\ = V \text{diag}(1/\sigma_1, \dots, 1/\sigma_m) U^T$$

largest

$$\text{So, } \kappa(A) = \sigma_1 / \sigma_m$$

We can **generalize** the definition of the **condition number** for a rectangular matrix $A \in \mathbb{R}^{m \times n}$ using the pseudo-inverse $A^\dagger$ and SVDs as

$$\kappa(A) := \|A\|_2 \cdot \|A^\dagger\|_2 \\ = \sigma_1 / \sigma_r \quad \begin{array}{l} r = \text{rank}(A) \\ \leq \min(m, n) \end{array}$$