


MAT 108

Winter 2021
LEFT BOARD



Elem \longrightarrow Adv math,
Reading and writing proofs,
Like essay - structure
- ideas

Eng is not precise enough.

So use predicate logic.

And some set theory,

Def: A proposition is a sentence which is true or false
 (T) (F)
 (or has a truth value).

Examples:

$$1+1=3$$

prop

Y

it is F

	Prop?
① This sentence is not a prop.	Y F
② I am liar.	N paradox

③ Lysol can kill viruses. 14

① If F then the sent is a prop. ✓
If T then not a prop x

① Call the sentence P:

If P is True then P is not a prop. } Is
so not T or F } T

If P is False then P is a prop. } Not
so either T or F } F

Using steps from Thm 1.1.1:

$$\sim P \wedge Q$$

by (h) $\sim (A \wedge B)$ is eq. $\sim A \vee \sim B$

~~by (a)~~ $\sim(\sim A)$ is eq. to A

$\sim P \wedge Q$ is eq. to $\sim[\sim(\sim P \wedge Q)]$
②

which is eq to

$$(b) \quad \sim [\sim(\sim P) \vee \sim Q]$$

which is eq to

$$(a) \quad \sim [P \vee \sim Q]$$

210106

Building new props from old:

$\sim Q$ ($\sim Q$) same.

$Q \wedge (P \vee R)$

Truth tables

P	Q	T	$\sim P$	$P \wedge Q$	$P \vee Q$	$P \Rightarrow Q$	$Q \Rightarrow P$	$(\sim Q) \Rightarrow (\sim P)$
T	T	T	F	T	T	T	T	T
T	F	T	F	F	T	F	T	F
F	T	T	T	F	T	T	F	T
F	F	T	T	F	F	T	T	T

equivalent

Note: $P \Rightarrow Q$ and $(\sim Q) \Rightarrow (\sim P)$ are equivalent. as is $\sim P \vee Q$
the contrapositive.

Eng

Ex:

P: My dog is hungry.

Q: My dog is inside.

$P \Rightarrow Q$: If my dog is hungry then it is inside.

$Q \Rightarrow P$: If my dog is inside then it is hungry.

$\sim Q \Rightarrow \sim P$: If my dog is outside then it is full.

$\sim P \vee Q$: Either my dog is full or it is inside.

1,3,4 have the same meaning.
2 is different.

$\sqrt{210108}$

Today quantifiers;
Complete. first order logic notation,

Ex:

Everyone I know likes chocolate



or dislikes coffee.

Rewriting this in logic:

$P(x)$ is x likes chocolate.

$Q(x)$ is x likes coffee.

Notation: A sentence like $P(x)$.
is an open sentence with variable x .

The above becomes:

$(\forall x \in \{\text{people I know}\}) P(x) \vee \sim Q(x)$

for all x in the set of people I know. \vee person I know.

or: In the universe (of discourse)

{people I know}.

$$(\forall x) P(x) \vee \sim Q(x)$$

Everyone I know who likes chocolate
also likes coffee.

equiv.

For everyone I know if they like chocolate
then they ~~also~~ like coffee

$$(\forall x \in \{\text{pp} \mid \text{I know}\}) P(x) \Rightarrow Q(x)$$

Compute $(\forall x \in \{pp \mid I \text{ know with } P(x) \text{ true}\})$
(equivalent) $(Q(x))$.
✓

Thm 1.3.1: If $P(x)$ is an open sentence then in any universe

(a) $\sim (\forall x) P(x)$ is eq. to $(\exists x) \sim P(x)$

(b) $\sim (\exists x) P(x)$ is eq. to $(\forall x) \sim P(x)$

Check example:

$$\sim (\forall x \in \{\text{ppl I know}\}) (P(x) \vee \sim Q(x)) \quad \text{eg } 6.4.1$$

$$(\exists x \in \{\dots\}) \sim (P(x) \vee \sim Q(x)) \quad \text{eg } 1.3.1$$

$$(\exists x \in \dots) (\sim P(x) \wedge Q(x)) \quad \text{eg } 6.4 \text{ DeMorgan}$$

Def: The truth set in a universe U
for an open sentence $P(x)$.
is all x in U for which $P(x)$
is true

26/11 Proofs:
see §1.7 pgs 64, 65, 66, 67

Examples:

Def: An integer $a \in \mathbb{Z}$ is

even if there is an integer n
with $a = 2n$.

$$(\exists n \in \mathbb{Z}) (a = 2n)$$

An integer $a \in \mathbb{Z}$ is
odd if $(\exists n \in \mathbb{Z}) (a = 2n + 1)$

Thm: If x is a real number
with $x^2 \leq 1$ then $x^2 - 7x > -10$

Proof: Assume x is a real number with
 $x^2 \leq 1$.

Hence $|x| \leq \sqrt{x^2} \leq \sqrt{1} = 1 < 2$.

Hence $x < 5$.

Hence $(x-2) < 0$ and $(x-5) < 0$

Then $(x-2)(x-5) \geq 0$.

Hence $x^2 - 7x + 10 \geq 0$.

Therefore $x^2 - 7x \geq -10$. q.e.d.

Approach:

$$x^2 + 6 > 5x$$

$$(x-2)(x-3) = x^2 - 5x + 6 > 0$$

$x > 3$
both pos

$x < 2$
both neg

\cap

$x < 1$

Proofs to grade? (bad example)

Thm: If a is an odd integer
then a^2+1 is an even integer.

"Proof": Let a . (2)

(1) [Then by squaring an odd
we get an odd.
(3) [An odd plus an odd is even.
So a^2+1 is even.

Problems: ① Why is this true?
② This is not a sentence.

should be

Let a be an odd integer.

③ is also not clear.

210113 Proof Structures;

Recall: § 1.7 64-67 should be reread.

From pg 67:

To start working out a proof consider:

Understand the statement

Logical form.

Assumptions and Conclusion.

Ideas

Step 3: Underline.

Try an example: eg $a=3$ and $b=6$ and
 $c=9$ claim is that 3 divides $6-9$
 $= -3$

Logic: $P \wedge Q \Rightarrow R$
If... then...

P is $a \text{ div } b$

Q is $a \text{ div } c$

R is $a \text{ div } b-c$

with quantifiers:
 $(\forall a, b, c \in \mathbb{Z})$
 $(\forall a \in \mathbb{Z})(\forall b \in \mathbb{Z})(\forall c \in \mathbb{Z})$

Ass & Concl:

Ass: P, Q

$$a \text{ div } b \quad \text{or} \quad (\exists n) \quad n \cdot a = b$$

Ass

$$a \text{ div } c \quad \text{or} \quad (\exists m) \quad m \cdot a = c$$

Ass

Concl: R

$$a \text{ div } b - c$$

Idea: With a, b, c as above

have

$$b - c = n \cdot a - m \cdot a =$$

last,

$$(n - m) \cdot a$$

this is an integer.

Contra position proof:

Claim: Assume that m is an integer.
If m^2 is an odd integer
then m is an odd integer.

Proof: Assume m is an integer.

Assume m is even.

Hence $\exists t \in \mathbb{Z}$ with $2 \cdot t = m$.

Hence $2t^2$ is also an integer.

Hence $m^2 = (2t)^2 = 4t^2 = 2(2t^2)$.

Hence m^2 is even.

There fore! If m^2 is odd then m is odd.
qed.

Example with cases and
proof by contradiction!

Idea of proof by contradiction:

To prove P , assume $\sim P$

and show Q and $\sim Q$.

U: understand.

eg $a=1$ $a^2 - 2 = -1$ ✓
 $a=2$ $a^2 - 2 = 2$ ✓
 $a=4$ $a^2 - 2 = 14$ ✓



L: Logic!

Right now: $(\forall a) P$

Plan: $(\forall a) \underbrace{\sim P \Rightarrow (Q \wedge \sim Q)}$

(to see these are equivalent:

$\swarrow \searrow$
 $P \vee \underbrace{(Q \wedge \sim Q)} \quad \text{or} \quad P$

Contradiction

$$\neg(Q \wedge \neg Q) \Rightarrow P$$

eg. ~~as~~ $\underline{\neg Q \vee Q} \Rightarrow P$

Ass: $\neg P$

Condi: Q and $\neg Q$

Ideas: $\neg P$ is $(4 \text{ divides } a^2 - 2)$

or $(\exists t) \text{ with } 4t = a^2 - 2$

Cases: \textcircled{a} a is even

⑥ a is odd

even: $a = 2s$ so $4t = (2s)^2 - 2 = 4s^2 - 2$

so $2t = 2s^2 - 1$

or $1 = 2s^2 - 2t = 2(s^2 - t)$

Q: 1 is even.

$\sim Q$ is clearly true.

Need to show Q

2/10/15 Mon. No Lect.
HW due Wed.

Recall: Claim ①: If a is an integer
and a^2 is even then a is even.
Pf: Earlier.

Claim ②: If a is an integer
then 4 does not divide $a^2 - 2$.

Proof: Note that 1 is not even.

Assume a is an integer and 4 divides $a^2 - 2$

Hence there is an int t with $4t = a^2 - 2$ ✓

$$\text{Hence } a^2 = 2(2t - 1)$$

so $2t - 1$
is an integer.

so a^2 is even

and by Claim ① a is even.

Hence there is an integer s with $a = 2s$

and $s^2 - t$ is an integer.

$$\text{Hence } 4t = (2s)^2 - 2 \quad \text{so } 1 = 2(s^2 - t) \text{ is even.}$$

Therefore 1 is even and 1 is not even
which is a contradiction. qed.

A: $\sim Q$ or p is even so $p = 2s$
 $\sim R$ or q is even so $q = 2k$] ①

Concl: $\sim p$ or q is not the smallest possible denominator.

I: $a = \frac{p}{q} = \frac{2s}{2k} = \frac{s}{k}$] ②
and $k < q$ and hence a smaller denom. so $\sim p$ ✓

Proof: Assume $a = \frac{p}{q}$ with
 p and q both even integers.

Hence there are integers s and k with
 $p = 2s$, $q = 2k$ so $a = \frac{p}{q} = \frac{2s}{2k} = \frac{s}{k}$

Therefore q is not the smallest possible
denominator q ed.

U! eg $1^2 + 15 = 16$
 $8 \cdot 1 = 8$ $16 \neq 8$ oops,
maybe such else

L! $(\exists n \in \mathbb{Z}) (n^2 + 15 < 8n)$,

A! No ass. work

Concl! $n^2 + 15 < 8n$

I! $n^2 + 15 - 8n < 0$

or $(n-5)(n-3) < 0$

need $n-5$ & $n-3$
to have different signs,

so take $n = 4$

210120) Pythagorean's Thm

Thm: $\sqrt{2}$ is irrational.

Plan Proof:

u: try $(\frac{7}{5})^2 = 1.96$

$$(\frac{10}{7})^2 = 2.040816\ldots$$

L: $\sim P$

P is $\sqrt{2}$ is rational

or (contradiction approach)

$$P \Rightarrow (Q \wedge \sim Q)$$

Ass: $\sqrt{2} = \frac{p}{q}$

Concl: \mathbb{Q} and $\sim \mathbb{Q}$.

still have not had to choose \mathbb{Q} .

I: If $\sqrt{2} = \frac{p}{q}$

③ ✓

Recall: If q is as small as possible

① Recall

Tidy.

p is odd or q is odd.

Compute $2 = \frac{p^2}{q^2}$ or $2q^2 = p^2$

④ ✓

so p^2 is even means:

also should be before this

② Recall.

p even.

q is odd.

and ④ $p = 2k$ and $q = 2m+1$ ⑤

Proof: Note that 1 is not even.

Say that if a is rational, $p \in \mathbb{Z}$ and $q \in \mathbb{N}$
with $a = \frac{p}{q}$ and q as small as possible
then $a = \frac{p}{q}$ is in reduced form.

Recall we proved last time that if $a = \frac{p}{q}$,
is a rat. number in reduced form then p or q
is odd.

Recall we proved before that if n is an int.
and n^2 is even then n is even.

For contradiction assume $\sqrt{2}$ is rational

and $\sqrt{2} = \frac{p}{q}$ is in reduced form.

Hence $2 = \frac{p^2}{q^2}$ so $2q^2 = p^2$ so p^2 is even so p

is even so $p = 2k$ for some int. k .

Hence q must be odd so $q = 2m+1$

for some int. m .

Hence $k^2 - 2m^2 - 2m$ is an integer and

$$4k^2 = (2k)^2 = p^2 = 2q^2 = 2(2m+1)^2 = 8m^2 + 8m + 2$$

Hence $2[k^2 - 2m^2 - 2m] = 1$ and 1 is even.

Therefore 1 is even and 1 is not even a contradiction
so $\sqrt{2}$ is irrational. qed

Recall $(\exists! x)(P(x))$

is equiv to:

$$(\exists x)(P(x)) \wedge (\forall u, v)[(P(u) \wedge P(v)) \Rightarrow (u=v)]$$

find an example.

Ans:

(A) (3) has more than 1] false

④ $(x-2)^2=0$

True

⑤ $\frac{4 \pm \sqrt{-4}}{2}$ not \mathbb{R}
no answers

7 false



③ Find a different example.
and done.

④ Example for \exists
More work for \forall .

Claim: $(\exists! x \in \mathbb{R}) \quad x^2 - 4x + 4 = 0$.

Proof! First show $(\exists x \in \mathbb{R}) \quad x^2 - 4x + 4 = 0$

by taking $x=2$ so $2^2 - 4 \cdot 2 + 4 = 0 \checkmark$.

Uniqueness: Assume $u^2 - 4u + 4 = 0$

and $v^2 - 4v + 4 = 0$,

Hence $(u-2)^2 = 0$

and $(v-2)^2 = 0$

so $u-2 = 0$

and $v-2 = 0$

so

$u=2$

and $v=2$ so $u=v$.



If P then Q .

~~or~~ or

$$P \Rightarrow Q$$

By cont:

eq.

$$[\neg (P \Rightarrow Q)] \Rightarrow (R \wedge \neg R)$$

~~Possibly can choose $R = \emptyset$~~

→ eq: $[\neg (\neg P \vee Q)] \Rightarrow (R \wedge \neg R)$

$$Q: (P \wedge \sim Q) \Rightarrow (R \wedge \sim R)$$

$$\rightarrow \left[\begin{array}{l} \cancel{Q: (P \wedge \sim Q) \Rightarrow (Q \wedge \sim Q)} \\ Q: P \wedge \sim Q \Rightarrow Q \end{array} \right]$$

or ^{maybe choose} $R \equiv P$

$$(P \wedge \sim Q) \Rightarrow (P \wedge \sim P)$$

$$\text{or } P \wedge \sim Q \Rightarrow \sim P$$

enough to show $\sim Q \Rightarrow \sim P$

210122 §1.8 Number Theory (for proofs).

Next week Set Theory (" "),

Recall: If a and b are integers
then a divides b iff there is an integer
 c with $a \cdot c = b$.

If p is an integer then p is prime
iff the only positive integers
dividing p are 1 and p .

Def (pg 77): If a, b and d are integers
nonzero

then $\gcd(a, b) = d$ iff

(i) d divides both a and b

(say d is a common divisor of a and b)

(ii) every common divisor of a and b
is at most d .

In first order logic:

$$(\forall a, b, d \in \mathbb{Z}_{\neq 0}) [\gcd(a, b) = d] \iff$$

$$[(\exists s, t \in \mathbb{Z}) (d \cdot s = a) \wedge (d \cdot t = b)]$$

$$\wedge [(\forall e \in \mathbb{Z}) [(\exists u, v \in \mathbb{Z}) (e \cdot u = a) \wedge (e \cdot v = b)] \Rightarrow (e \leq d)]$$

d is a common
div.

e is a common
div of a & b

Brk Rm: Translate lcm def.
into logic.

Claims: $\{10, 11, 12, \dots\}$

- (a) $(\forall a \in \mathbb{N}_{\geq 10}) (\exists b \in \mathbb{N}) (gcd(a, b) = 1) \wedge (a \leq b)$
- (b) $(\exists a \in \mathbb{N}_{\geq 10}) (\forall b \in \mathbb{N}) (gcd(a, b) = 1) \vee (a \leq b)$

Proof sketch for a:

Assume $a \geq 10$ is an integer.

Choose $b = a + 1$.

Note that 1 and -1 are the only
divisors of 1.

Assume that d is a common div. of a

and $b = a + 1$.

Hence there are integers s and t .

with $d \cdot t = a + 1$ and $d \cdot s = a$

so $d \cdot (t - s) = a + 1 - a = 1$

so d is 1 or -1 .

Hence $(a, d) = 1$.

qed.

2/10/25 More number theory from §1.8.

Division and Euclid's Alg.

Def; If a, b, x, y, n are integers

and $n = a \cdot x + b \cdot y$ then

n is a linear combination of a and b

Thm 1.8.1 (Prove later by induction)

If a and b are nonzero integers.

then $\gcd(a, b)$ is equal to the smallest

positive linear combin. of a and b .

Thm 2.5.1 (Division Alg).

If a and b are nonzero integers.

there is a unique pair of integers

q and r with

$$b = a \cdot q + r$$

and $0 \leq r < |a|$.

Notation for Euclid's Alg:

$$b = a \cdot q_1 + r_1$$

$$a = r_1 \cdot q_2 + r_2$$

$$r_1 = r_2 \cdot q_3 + r_3$$

...

$$r_{k-2} = r_{k-1} \cdot q_k + r_k$$

$$r_{k-1} = r_k \cdot q_{k+1}$$

Thm: 1.8.2: If $b \geq a > 0$ are integers then $\gcd(a, b) = r_k$ from Euclid's algorithm.

Brk. Rm: Apply Euclid's Alg to $b = 256 \geq a = 81 > 0$

Find $r_k = 1$ and k , and the q_i 's.
 $k = 3$ $q_1 = 13$ $q_2 = 3$ $q_3 = 1$

Claim: (1.8.3):

If a, b and p are integers with p prime
and p divides ab then p divides a
or p divides b .

Proof: Assume a, b and p are integers with
 p prime, p dividing ab but not a .

Hence the only pos. divisors of p are

1 and p so $\gcd(p, a) = 1$.

Hence by Thm 1.8.1 there are integers x

and y with $1 = x \cdot p + y \cdot a$.

Also there is an int. n with
 $ab = np$ and $b = b \cdot x \cdot p + y \cdot a - b$
 so $b = b \cdot x \cdot p + y \cdot n \cdot p$
 $= [bx + yn] \cdot p$ and $bx + yn$ is
 an integer.

Therefore p divides b . q.e.d.

Set Notation examples:

$$\{3, 4, 5, 6\} = \{x \in \mathbb{Z} \mid x \geq 3, x \leq 6\}$$

$$= \{ \text{" " : " " } \}$$

$$= \{x \mid x \in \mathbb{Z}, 3 \leq x \leq 6\}$$

has 4 elements.

Write $3 \in \{3, 4, 5, 6\}$

\nwarrow is an element of

$2 \notin \{3, 4, 5, 6\}$ \nwarrow is not " " " "

$$\{3, 4\} \subseteq \{3, 4, 5, 6\}$$

\nwarrow is a subset of

$$\{3\} \not\subseteq \{3, 4, 5, 6\}$$

3 \notin $\{3, 4, 5, 6\}$

$\{3\} \in \{\{3\}, \{1, 5, 6\}\}$ has 2 elts

Power sets;

$$P(\{2, 3\}) = \{\{\}, \{2\}, \{3\}, \{2, 3\}\}$$

The set of subsets.

$\varnothing =$ the empty set

Which are true:

$\exists A, B, C$ sets with

⑥ $B \subsetneq A$, $B \subseteq C$, ~~$A \subseteq C$~~ and $C \not\subseteq A$

True

subset but not equal to } proper subset.

① True $A \subseteq B$, $B \not\subseteq C$, $A \subseteq C$

② $A \subseteq B$, $B \subseteq C$, $C \subseteq A$

True $A=B=C$

③ $A \subsetneq B$, $B \subsetneq C$, $C \subsetneq A$

False

④ $B \subseteq A$, $B \not\subseteq C$, $A \subseteq C$

False

210129 Set operations.

Related to operations on predicates.

Notation: If A and B are sets write

① $A \cup B = \{x \mid (x \in A) \vee (x \in B)\}$

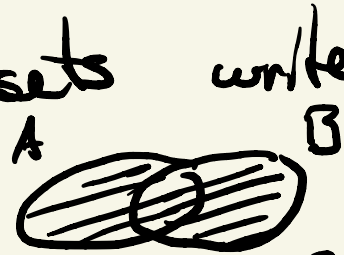
② $A \cap B = \{x \mid (x \in A) \wedge (x \in B)\}$

③ $A - B = \{x \mid (x \in A) \wedge (x \notin B)\}$

union

intersection

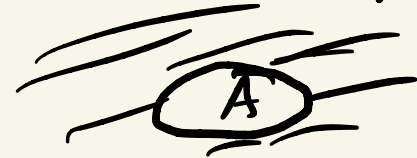
difference



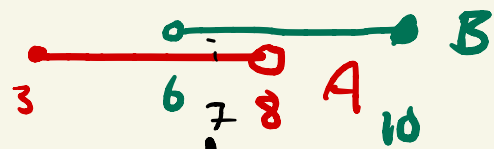
If A is a subset of a universe U write:

④ $A^c = U - A$

complement



$$= \{x \in U \mid x \notin A\}$$



Example: $A = [3, 8) \subseteq \mathbb{R} = U$
 $B = (6, 10] \subseteq \mathbb{R}$

- Find
- ① $A \cup B$
 - ② $A \cap B$
 - ③ $A - B$
 - ④ A^c
 - ⑤ $A \cap B^c$

- Ans:
- ① $[3, 10]$
 - ② $(6, 8)$
 - ③ $[3, 6]$
 - ④ $(-\infty, 3) \cup [8, \infty)$
 - ⑤ this just $A - B$ which is $[3, 6]$.

Dictionary: If $P(x)$ and $Q(x)$ are open propositions with variable x in U a universe

take $A = \text{Truth}(P) = \{x \in U \mid P(x) \text{ is true}\}$

$B = \text{Truth}(Q) = \{x \in U \mid Q(x) \text{ is true}\},$

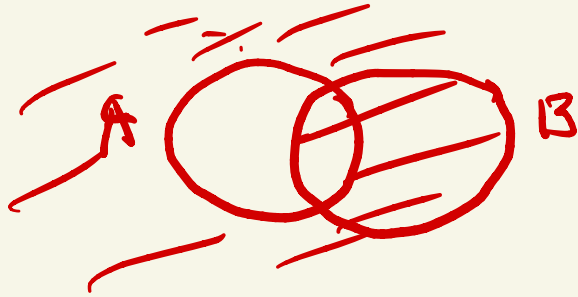
$$\text{Truth}(P \wedge Q) = A \cap B$$

$$\text{Truth}(P \vee Q) = A \cup B$$

$$\text{Truth}(\sim P) = A^c$$

$$\text{Truth}(P \wedge \sim Q) = A - B$$

$$\text{Truth } (P \Rightarrow Q) = A^c \cup B$$



$$= (A - B)^c$$

Truth Table

simila.

Venn Daiagram

210201 More set notation:

Products then Families:

Def: If A and B are sets then

$$A \times B = \{(a, b) \mid a \in A \text{ and } b \in B\}.$$

the product set

$$\text{Ex: } \{1, 2\} \times \{2, 3\} = \{(1, 2), (1, 3), (2, 2), (2, 3)\}$$

Ex: Which one is true?

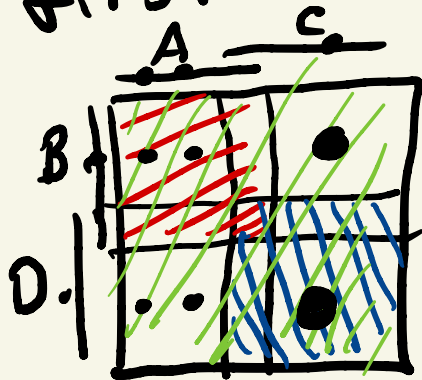
Prove one and find a counterexample.

So the other.

$(\forall A, B, C, D \text{ sets}) (A \times B) \cup (C \times D) \subseteq (A \cup C) \times (B \cup D)$
" " " " "

Ans: The first is true.

U:



$A \times B$

$C \times D$

$(A \cup C) \times (B \cup D)$

For a counterexample to \supseteq

L: If $\underbrace{P}_{\text{P}} \underbrace{(x,y) \in (A \times B)}_{\text{Q}} \text{ or } \underbrace{(x,y) \in (C \times D)}_{\text{Q}}$
 then $\underbrace{(x,y) \in (A \cup C) \times (B \cup D)}_{\text{R}}$.

C: Assume: P and conclude R.
 and Assm Q and conclude R.

Sets or families or collections of sets.

Ex: ①

$$Q = \left\{ \underbrace{\{1, 2, 3\}}_{A_a}, \underbrace{\{2, 3, 4\}}_{A_b}, \underbrace{\{3, 4, 5\}}_{A_c} \right\} = \{A_a \mid a \in \Delta\}$$

$$\Delta = \{a, b, c\} \text{] indexing set.}$$

$$\textcircled{2} \quad \mathcal{B} = \{ [x, x+3] \mid 0 \leq x < 2 \}$$

$$= \{ B_\alpha \mid \alpha \in \Delta \}$$

$$\text{if } B_\alpha = [\alpha, \alpha+3]$$

$$\text{and } \Delta = [0, 2) \text{] indexing set}$$

Ex ①

$$\bigcap_{A \in \mathcal{Q}} A = \bigcap_{\alpha \in \{a, b, c\}} A = \{3\}$$

②

$$\bigcap_{B \in \mathcal{B}} B = \bigcap_{\alpha \in [0, 2)} B_\alpha = [2, 3]$$

210203 Last proof technique: Induction

Example:

Claim: For every $n \in \mathbb{N} = \{1, 2, 3, \dots\}$
it is true that
$$n^2 = 1 + 3 + 5 + \dots + (2n-1)$$

Proof: Check the base case of $n=1$
which is $1^2 = 1$ which is true.

Assume for induction that

$$n^2 = 1 + 3 + 5 + \dots + (2n-1)$$

Hence $(n+1)^2 = n^2 + 2n + 1$

$$= [1 + 3 + \dots + (2n-1)] + [2n+1]$$

$$= 1 + 3 + \dots + [2(n+1)-1]$$

Therefore by PMI the claim holds.
 qed.

principle of
mathematical
induction

could add [since $2n+1$ is the odd number
following $2n-1$.

Proof Plan:

W:	$n=1$	$1+3 < 5$	✓
	$n=3$	$3+3 < 5$	✓

L: $(\forall n \in \mathbb{N}) P(n)$

use induction.

C: • Prove $P(1)$ (base case).

• Assume $P(n)$ } induction

show $P(n+1)$ 1 step,

I: base case $n=1$ see above.

ind. step:

Assume: $n+3 < 5n^2$

Check: $(n+1)+3 = (n+3)+1 < 5n^2+1$

$$5n^2+1 < 5n^2+10n+5 = 5(n+1)^2$$

Generalized Prin. of Math. Ind:

Show $(\forall n \in \mathbb{N}) \quad P(n)$.

by:

- Show $P(1)$ (base case).
- Assume for every $n \leq m$
have $P(n)$.

Show $P(m+1)$.

Example:

Claim: $\sqrt{2}$ is irrational.

210205 Midterm next Wednesday;
On web is an old exam;
Covers Ch1 & Ch2 (except 26).

More induction proofs;

Def: (Fibonacci numbers):

Inductive definition:

$$f_1 = 1, f_2 = 1 \text{ and } f_{n+2} = f_{n+1} + f_n \\ \text{if } n > 0.$$

$$\text{Ex: } f_3 = 2, f_4 = 3, f_5 = 5, f_6 = 8, f_7 = 13, \dots$$

Claim ①: If f_n are the Fibonacci numbers
then for every $n \geq 1$ have $\gcd(f_n, f_{n+1}) = 1$

Claim ②: If f_n are the Fibonacci numbers
then for every $n \geq 1$ have that 3 divides
 f_{4n} .

?
∴ follows from:

Call
This a Lemma
and prove it.

$$\gcd(b, a) = \gcd(b-a, a)$$

$f_{n+1} + f_n$ f_{n+1}

Focus on the Lemma:

①

Check: $\{d \mid d \mid b \text{ and } d \mid a\}$
 $= \{e \mid e \mid b-a \text{ and } e \mid a\}.$

check: \subseteq : If $d \mid b$ and $d \mid a$
want to show $d \mid b-a$ and $d \mid a$
but if $b = sd$ and $a = td$
then $b-a = sd - td = (s-t)d.$

So $d \mid b-a$ ✓.

\supseteq ; similar.

Claim: If f_n are the Fibonacci numbers then for every natural number n we have

$$\gcd(f_n, f_{n+1}) = 1.$$

Proof: Use the following Lemma:

Lemma: If $b \geq a > 0$ then $\gcd(b, a) = \gcd(b-a, a)$.

Proof of Lemma: It suffices to show that that b and a have the same common divisors as $b-a$ and a .

If d divides b and a then $b = sd$ and
 $a = td$ so $b - a = (s - t)d$.

If d divides $b - a$ and a then $b - a = ud$
and $a = vd$ so $b = (b - a) + a = (u + v)d$.
hence the lemma holds.

Prove the claim by induction.

For the $n=1$ base case need $\gcd(h_1) = 1$
which is true,

For induction assume $\gcd(f_n, f_{n+1}) = 1$

Hence $\gcd(f_{n+2}, f_{n+1}) = \gcd(f_{n+1} + f_n, f_{n+1})$.

Hence by the lemma $\gcd(f_{n+2}, f_{n+1}) = \gcd(f_n, f_{n+1})$

which is 1 by the ind. hyp.
qed.

Claim ②: If f_n are the Fibonacci numbers then every f_{4n} is divis. by 3.

Proof: For induction consider the base case $n=1$ which is $f_4 = 3$ is div. by 3.

For ind. assume f_{4n} is div. by 3.

$$\begin{aligned}
 \text{Hence } f_{4(n+1)} &= f_{4n+3} + f_{4n+2} \\
 &= f_{4n+1} + 2f_{4n+2} \\
 &= 3f_{4n+1} + 2f_{4n}
 \end{aligned}$$

which is divisible by 3 since f_{4n} is.
 Therefore by PMI. the claim holds. end.

$$f_{4n} = 3 \cdot s \quad \text{so}$$

$$f_{4(n+1)} = 3 \cdot f_{4n+1} + 2 \cdot 3 \cdot s = 3 [f_{4n+1} + 2s]$$

Horses are all the same color:

Case 1

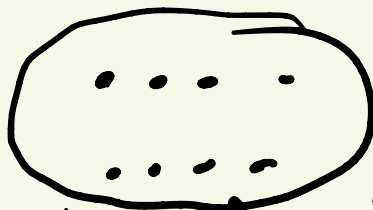
• ✓

color. $\begin{bmatrix} 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$ same color

Problem!

The induction proof is
wrong for case $n+1=2$

Example: Nim:



move; remove any number of coins from a pile.

Start with 2 equal piles;

2nd player has a winning strategy.

win: take last coin.

Proof: by Gen. ind.

Base case: one coin in each pile ✓

For ind. assume there is a
winning 2nd ply strat. for any
number of coins at most n
in 2 equal piles -

Show: if there are 2 piles with
 $n+1$ each. then after ply
takes r from one pile.

The second can take r from both.
and by the ind. hyp. there is a winning
strat. for 2 piles of $(n+1-r)$.

GI Ind: Assume $\forall k \leq n$ have $P(k)$.
Prove $P(n+1)$

210208 | Wednesday ; Midterm

1 sheet notes, (both sides),

Chapt 1, 2 (except 2.6).

Regular zoom location.

Instructions up this afternoon,

Cameras: Point at your work.

Try out during quiz tomorrow.

Uploading similar to Homework

Arrive ~10 min early to check IP...

Exam: Both on screen
and via chat box link.

50 min exam.

Old exam and practice exam on
web page (answers today).

Upload after 50 min.

gcd proofs
set proofs.
families proofs.

induction.

Proof Sketch:

Und:



$$\bigcup_{\alpha \in \Gamma} A_\alpha$$

$$\bigcup_{\alpha \in \Delta} A_\alpha$$



Logic: $(a \in \bigcup_{\alpha \in \Gamma} A_\alpha) \Rightarrow (a \in \bigcup_{\alpha \in \Delta} A_\alpha)$.

Ass/ccl: Assume $a \in \bigcup_{\alpha \in \Gamma} A_\alpha$

or $(\exists \alpha \in \Gamma) (a \in A_\alpha)$.

$$\left. \begin{array}{l} \bigcup_{\alpha \in \Delta} \leftrightarrow \exists \alpha \\ \bigcap_{\alpha \in \Delta} \leftrightarrow \forall \alpha \end{array} \right\}$$

2.4.2 Concl: $a \in \bigcup_{\alpha \in \Delta} A_\alpha$

or $(\exists \alpha \in \Delta) (a \in A_\alpha)$.

Idea: If $a \in \Gamma$ and $a \in A_\alpha$
then since $\Gamma \subseteq \Delta$ also $a \in \Delta$,
as needed.

Proof sketch:

Und!, eg $m=5, n=13$

$$13 = \underbrace{2 \cdot 5}_q + \underbrace{3}_r$$

Logic! Want $\exists!$

which requires ① $\exists r, q$ with ---

② any 2 solns are the same (!)

$$(\forall n \in \mathbb{N}) (\forall m, 0 < m \leq n) \quad P(n) \\ (\exists! q, r) [(n = q \cdot m + r) \wedge (0 \leq r < m)]$$

Might use induction: $(\forall n \in \mathbb{N}) P(n)$

Ass/Concl: Assume $P(m)$ for every $m < n$,
Conclude $P(n)$.

Generalized
induction.

Ideas: Uniqueness:

Assume there are two solns
Show be same:

$$n = q \cdot m + r$$

$$\text{or } n = \tilde{q} \cdot m + \tilde{r}$$

$$0 \leq r < m$$

$$0 \leq \tilde{r} < m$$

$$\begin{aligned} \text{Hence; } 0 = n - n &= q \cdot m + r - \tilde{q} \cdot m - \tilde{r} \\ &= (q - \tilde{q}) \cdot m + (r - \tilde{r}). \end{aligned}$$

$$\tilde{r} - r = (q - \tilde{q}) \cdot m$$

so m divides $\tilde{r} - r$

so

Also

$$-m < \tilde{r} - r < m$$

so

$$(\tilde{r} - r) = m \cdot s$$

$$-m < m \cdot s < m$$

$$-1 < s < 1$$

Hence

$$\tilde{r} - r = 0$$

so

$$\tilde{r} = r.$$

$$\text{so } 0 = \tilde{r} - r = (q - \tilde{q}) \cdot m$$

$$\text{so } 0 = q - \tilde{q} \quad \text{so } q = \tilde{q}$$

210212 Set Theory and the rest of math!

300 BC on (Euclid), Geometry
as foundation.

1800's shift to set theory.

Next:

Relations \rightarrow functions
3.1 4

Set Products.
(2.2)

\downarrow
Equivalences \rightarrow Modular Arithmetic
3.2 Related to gcd.
3.4

Example: $A = \{1, 2, 3, 4\}$

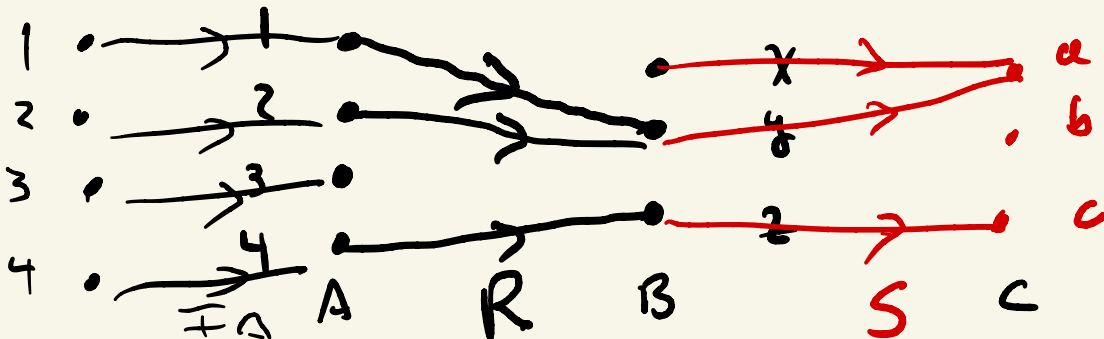
$B = \{x, y, z\}$

$R = \{(1, y), (2, y), (4, z)\}$

just
one
example

eg: $2Ry$ and $2Rx$

Digraph associated to R (like truth table or Venn diagram)

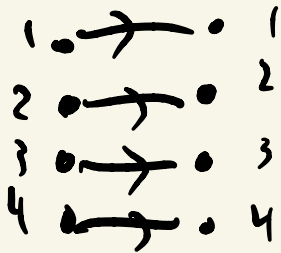


eg: $\text{Dom } (R) = \{1, 2, 4\}$

$\text{Rng } (R) = \{y, z\}$

Ex: $I_{\{1,2,3,4\}} = \{(1,1), (2,2), (3,3), (4,4)\}$.

write
digraph:



$R^{-1} = \{(y, 1), (y, 2), (z, 4)\}$

Ex: R as above and S a reln.
from B to $C = \{a, b, c\}$.

eg $S = \{(x, a), (y, a), (z, c)\}$
then $S \circ R = \{(2, a), (1, a), (4, c)\}$

Claim: If A and B are sets and
 R is a relation from A to B then
 $I_B \circ R = R$. Here I_B is the
identity relation on B .

Proof: Assume $(a, b) \in I_B \circ R$.

Hence there is $\tilde{b} \in B$ with $a R \tilde{b}$ and $\tilde{b} I_B b$. Hence $b = \tilde{b}$ and $a R b$.

Hence $(a, b) \in R$.

Assume $(a, b) \in R$. Hence if $\tilde{b} = b$
then $a R \tilde{b}$ and $\tilde{b} I_B b$ so $(a, b) \in I_B \circ R$.

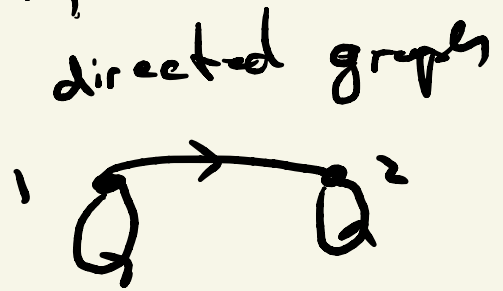
QED

210217 Equivalence Relations.

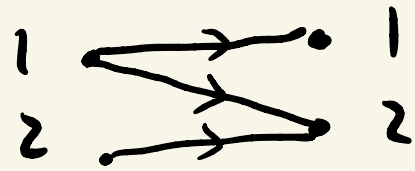
Recall: A relation from A to B
is a subset $R \subseteq A \times B$.

Example: $I_A = \{(x, x) \mid x \in A\}$

Ex: $R = \{(1, 1), (1, 2), (2, 2)\}$



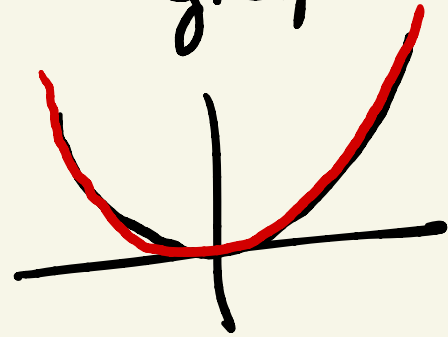
for a relation
on $\{1, 2\}$.



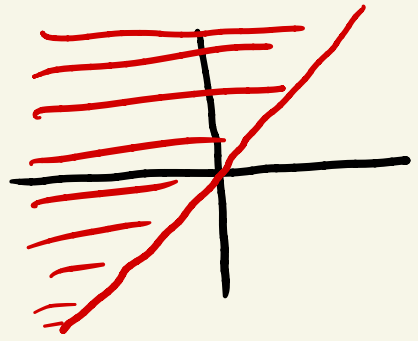
same reln R

Ex: Relations on \mathbb{R} = real numbers,
graph

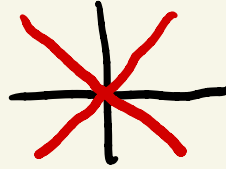
$$T = \{(x, y) \mid x^2 = y\}$$



$$U = \{(x, y) \mid x \leq y\}$$



$$V = \{(x, y) \mid x = y \text{ or } x = -y\}$$



Sketch:

Logic:

$$R = \{(x, y) \mid x R y\}.$$

$$R^{-1} = \{(x, y) \mid y R x\},$$

$\text{Dom}(R^{-1})$ and $\text{Rng}(R)$ are

sets so \subseteq are \supseteq .

try together:

$$\begin{aligned} \text{Rng}(R) &= \{y \mid (\exists x \in A) \ x R y\} \\ &= \{y \mid (\exists x \in A) \ y R^{-1} x\} \end{aligned}$$

$$\parallel$$

$$\text{Dom}(R^{-1}) = \{y \mid (\exists x \in A) \ y R^{-1} x\}$$

Properties of some relations;

Def: If R is a relation on a set A .

- ① R is reflexive if $(\forall x \in A) \quad x R x$
- ② R is symmetric if $(\forall x, y \in A) \quad x R y \Rightarrow y R x$
- ③ R is transitive if $(\forall x, y, z \in A)$
 $(x R y) \wedge (y R z) \Rightarrow x R z$
- ④ R is an equivalence relation if
it is reflexive, symmetric and transitive.

Claim: A relation R on A is symmetric iff $R = R^{-1}$.

210219 Proof: Assume R is a symmetric relation on A and $(a, b) \in R$. Hence $a R b$ so $b R a$ by symmetry of R so $a R^{-1} b$ so $(a, b) \in R^{-1}$.

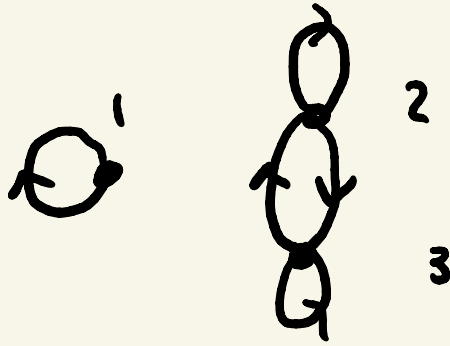
Therefore $R \subseteq R^{-1}$.
Assume R is sym on A and $(a, b) \in R^{-1}$.

Hence $a R^{-1} b$ so $b R a$ so $a R b$
so $(a, b) \in R$. Therefore $R^{-1} = R$.

For \Leftarrow assume $R = R^{-1}$ and
 $a R b$. Hence $(a, b) \in R = R^{-1}$ so
 $(a, b) \in R^{-1}$ so $a R^{-1} b$ so $b R a$.
Therefore R is symmetric. qed.

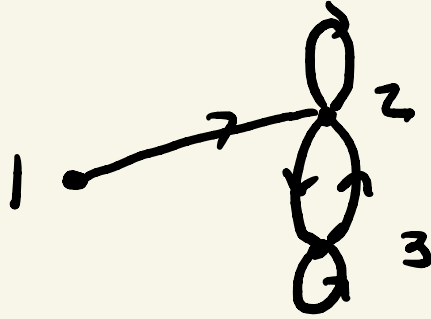
Examples:

$R =$
(equiv)



a relation on $A = \{1, 2, 3\}$

$S =$
(not equiv)



$$A/R = \left\{ \underbrace{\{1\}}_1, \underbrace{\{2, 3\}}_2 \underbrace{\quad}_3 \right\}$$

$$A/S = \{ \overset{''}{\{2\}}, \overset{''}{\{2, 3\}} \}$$

$$V = \{ (x, y) \in \mathbb{Z}^2 \mid x^2 = y^2 \}, \text{ a reln. on } \mathbb{Z}.$$



$$\mathbb{Z}/\sim = \{ \{0\}, \{1, -1\}, \{2, -2\}, \dots \}$$

$$= \{ \overset{''}{\{n, -n\}} \mid n \in \mathbb{Z} \}$$

$$= \{\{0\}\} \cup \{\{n, -n\} \mid n \in \mathbb{N}\}$$

$$U = \{(x, y) \in \mathbb{Z}^2 \mid 3 \text{ divides } x - y\}$$

a reln on ~~\mathbb{Z}~~

$$\mathbb{Z}/U = \left\{ \begin{array}{l} \{3, 0, 6, 9, 12, -3, \dots\}, \\ \{1, 4, 7, 10, -2, -5, \dots\}, \\ \{2, 5, 8, -1, -4, \dots\} \end{array} \right\}$$

$\begin{array}{l} \parallel \\ 3U3 \\ 3U0 \\ 3U6 \\ 3U9 \\ 3U-3 \end{array} \quad \begin{array}{l} \parallel \\ \frac{1}{3} \\ \parallel \\ \frac{1}{3} \\ \parallel \\ 0 \\ \parallel \\ \frac{1}{6} \end{array} \quad \begin{array}{l} \parallel \\ -5 \\ \parallel \\ \frac{1}{1} \\ \parallel \\ \frac{1}{4} \\ \parallel \\ -2 \end{array}$

$$\mathbb{Z}/u = \{ \overline{0}, \overline{1}, \overline{2} \}$$

$$= \{ \{3n \mid n \in \mathbb{Z}\}, \{3n+1 \mid n \in \mathbb{Z}\}, \{3n+2 \mid n \in \mathbb{Z}\} \}$$

$$= \{ \{3n+s \mid n \in \mathbb{Z}\} \mid s \in \{0,1,2\} \}$$

Note: U and V are both equiv. relns.

Proof sketch:

Un: See above examples.

Logic: ~~IF~~ $(R \text{ is eq.}) \Rightarrow (A/R \text{ a put}).$
 $(P_{ref} \wedge P_{sym} \wedge P_{trans}) \Rightarrow (Q_i \wedge Q_{ii} \wedge Q_{iii})$
 Q_i, Q_{ii}, Q_{iii}

I: $Q_i: \phi \notin A/R.$

if $\bar{x} \in A/R$ then: need $\bar{x} \neq \phi.$
but $x R x$ since $R \text{ refl.}$
so $x \in \bar{x} \neq \phi.$

Qiii: $\left[\begin{array}{l} \text{If } x \in A \text{ then} \\ x \in \bigcup_{\bar{y} \in \overline{F}} \bar{y} \end{array} \right. \text{ since } x \in \overline{x}$

~~not~~

Qii: $\forall \bar{x}, \bar{y} \in A/R \text{ have } (\bar{x} = \bar{y}) \vee$
 $(\bar{x} \cap \bar{y} = \emptyset)$

equiv: $(\bar{x} \cap \bar{y} \neq \emptyset) \Rightarrow (\bar{x} = \bar{y})$

Assume $z \in \bar{x} \cap \bar{y}$
 show $\bar{x} \subseteq \bar{y}$ (also need $\bar{x} \supseteq \bar{y}$).
 so assume $u \in \bar{x}$ and show $u \in \bar{y}$.

Assume $z \in \bar{x}, z \in \bar{y}, u \in \bar{x}$ and show $u \in \bar{y}$.
 $\Downarrow \quad \Downarrow \quad \Downarrow \quad \Uparrow$
~~Rec~~ $x R z \quad y R z \quad x R u \quad y R u$

by sym: $z R x$

so by trans here $y R x$ and $y R u$
so done and $\bar{x} \leq \bar{y}$.

210222 Next § 3.4: Modular Arithmetic.

Proof of Thm 3.3.1

Assume R is an equivalence relation
on a nonempty set A .

If $x \in A$ then $\bar{x} \in A/R$ and since R
is reflexive have $x R x$ so $x \in \bar{x}$.

Hence if $x \in A$ then $x \in \bar{x} \subseteq \bigcup_{\bar{y} \in A/R} \bar{y}$ hence

$$\bigcup_{\bar{y} \in A/R} \bar{y} = A.$$

Also if $\bar{x} \in A/R$ then $x \in \bar{x}$ so $\bar{x} \neq \emptyset$.

If $z \in \bar{y} \cap \bar{x}$ and $w \in \bar{x}$

then $y R z$, $x R z$ and $x R w$

so by symmetry $z R x$ and using transitivity
twice have $y R x$ and $y R w$

so $w \in \bar{y}$. Therefore $\bar{x} \subseteq \bar{y}$.

Similarly $\bar{y} \subseteq \bar{x}$ so $\bar{y} = \bar{x}$. qed.

Examples: In \mathbb{Z}_7 :

Find: or find the remainder after div by 7 of:

(a) $3 + 5 = \bar{8} = \bar{1}$ Ansi

(1)

(b) $3 - 5 = \bar{15} = \bar{1}$

(1)

(c) $5^3 = (-2)^3 = \bar{-8} = \bar{-1} = \bar{6}$

(6)

(d) $215 + 698 = \bar{5} + \bar{-2} = \bar{3}$

(3)

(e) $215 \cdot 698 = \bar{5} \cdot \bar{-2} = \bar{-10} = \bar{4}$

(4)

⑤

$$215^{698} \Rightarrow \overline{5}^{698} = \overline{(-2)}^{698} = \overline{(-2)}^{6(116)+2} = 1^{116} \cdot \overline{(-2)}^2 = \overline{4}$$

$$\overline{(-2)}^2 = \overline{4}$$

$$\overline{(-2)}^3 = \overline{-8} = \overline{-1}$$

$$\overline{(-2)}^6 = \overline{(-1)} \cdot \overline{(-1)} = \overline{1}$$

④

Try the same ①, ②, ⑤ in \mathbb{Z}_9

Ans: ① $\overline{8} + \overline{5} = \overline{13} = \overline{4}$

② $\overline{-1} \cdot \overline{5} = \overline{-5} = \overline{4}$

⑤ $\overline{(-1)}^{698} = \overline{1}$

④

④

①

RSA is the main public key
cryptosystem,

encrypt:

publish

698

~~encoder:~~

data

$215 \rightarrow 215^{698}$
" \oplus

hide

5

decode: $4^5 \rightarrow$

Sketch: Examples above: $m=7$

$$8 = 1 \quad (\text{mod } 7)$$

$$-2 = 5 \quad (\text{mod } 7)$$

$$6 = 6 \quad (\text{mod } 7).$$

Logic: $P \wedge Q \Rightarrow R.$

Assume P, Q show $R.$

$a \equiv_m c$ or $a-c$ is div. by $m.$

or there is k with $mk = a-c$

and $m_{\underline{l}} = b - d$

so $m_k + m_{\underline{l}} = (a - c) + (b - d)$
 $= (a + b) - (c + d)$

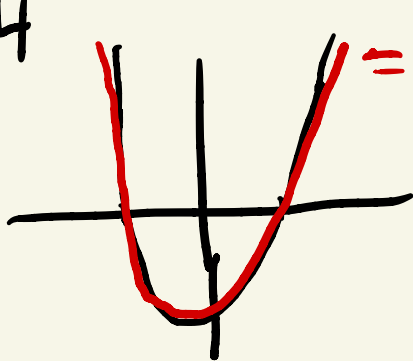
210224

Function: in a set theory framework.

Ex: formula for a function from \mathbb{R} to \mathbb{R} .

① $f(x) = x^2 - 4$

graph:



$= \text{graph}(f)$

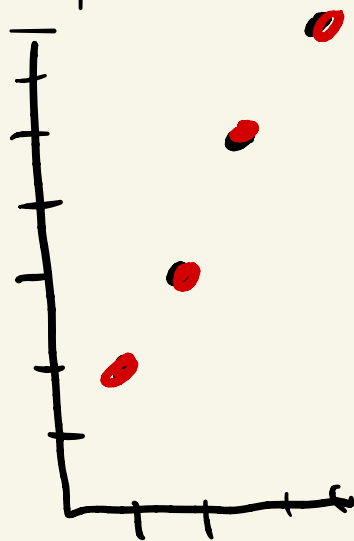
$= \{(x, y) \in \mathbb{R}^2 \mid y = x^2 - 4\}$

Ex: A sequence is a function from \mathbb{N}

② to some set.

eg: from \mathbb{N} to \mathbb{N} might be:

$p(n)$ = the n^{th} prime number.
 (so $p(1)=2$, $p(2)=3$, $p(3)=5$...).



$$= \text{graph} = \{ (n, m) \in \mathbb{N}^2 \mid m \text{ is the } n^{\text{th}} \text{ prime} \}$$

Vertical line test:

The graph of a relation is the graph of a function if it meets every vertical

line in exactly one point.

Proof: Assume $x \in \mathbb{R}$. Hence $y = x^2 - 4 \in \mathbb{R}$.
and $f(x) = y$ so $\text{domain}(f) = \mathbb{R}$.

Assume $f(x) = y$ and $f(x) = z$.

Hence $y = x^2 - 4 = z$. qed.

Claim: The inverse to ① is not a function.
 $g = \{(y, x) \in \mathbb{R}^2 \mid y = x^2 - 4\}.$

Sketch: either show $(-12, *)$ has no
solns so $-12 \notin \text{Dom}(g)$.

or show there are two solns.
 $(0, 2)$ & $(0, -2)$,
looks easier.

Proof: $(0, 2)$ and $(0, -2)$ are both in g .
(since $0 = 2^2 - 4$ and $0 = (-2)^2 - 4$) qed.

Claim: If f is a function from A to A and f is an equivalence relation on A then $f = I_A$.

Proof: Assume A is a set, f is a function from A to A and f is an equivalence relation.

If $a \in A$ then since f is reflexive $(a, a) \in f$ so $I_A \subseteq f$.

If $(b, c) \in f$ with $b \neq c$

then $(b, b) \in f$ also so
 by property (ii) of functions.
 $c = b$ a contradiction

hence $f \in I_A$.

qed.

Examples:

$$\textcircled{1} \mathbb{Z}_3 \xrightarrow{f} \mathbb{Z}_6$$

$$f(\bar{x}) = [2x]$$

$$\bar{x} \in \mathbb{Z}_3 \quad \{\bar{0}, \bar{1}, \bar{2}\}$$

$$[x] \in \mathbb{Z}_6 \quad \{[0], [1], [2], [3], [4], [5]\}$$

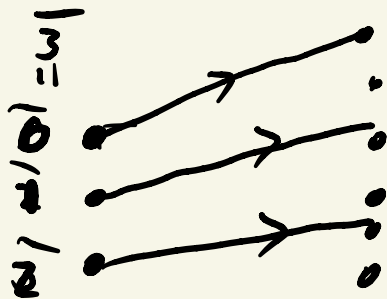
$$f = \{(\bar{x}, [2x]) \mid x \in \mathbb{Z}\}$$

$$\textcircled{2} \quad \mathbb{Z}_3 \xrightarrow{g} \mathbb{Z}_6$$

$$g = \{(\bar{x}, [x]) \mid x \in \mathbb{Z}\}.$$

$$g(\bar{x}) = [x]$$

①



$$[0] = [6]$$

$$[1]$$

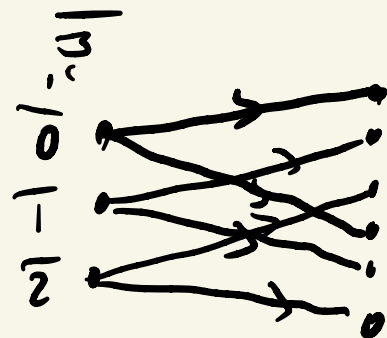
$$[2]$$

$$[3]$$

$$[4]$$

$$[5]$$

②



$$[0]$$

$$[3]$$

210226

write

$$\bar{x} \in \mathbb{Z}_3 = \{\bar{0}, \bar{1}, \bar{2}\}$$

" $\frac{1}{3}$...

$$[x] \in \mathbb{Z}_6 = \{[0], [1], [2], [3], [4], [5]\}$$

" $\{6\}$...

Claim: The relation $f(\bar{x}) = [x]$
 is not a well defined function
 from \mathbb{Z}_3 to \mathbb{Z}_6 .

Proof: From the definition
 $(\bar{0}, [0])$ and $(\bar{3}, [3])$ are in f
but $\bar{0} = \bar{3}$ and $[0] \neq [3]$. qed

Claim: The relation $g([x]) = \bar{x}$
is a (well defined) function from
 \mathbb{Z}_6 to \mathbb{Z}_3 .

Proof:

(i) If $[n] \in \mathbb{Z}_6$ then $([n], \bar{n}) \in g$.

(ii) If $([k], \bar{k})$ and $([l], \bar{l})$ are in g

and $[k] = [l]$ then there is t with

$$k - l = 6t = 3 \cdot 2t$$

so $\overline{k} = \overline{l}$. q.e.d.

Inverse function:

If f is a function from A to B
then sometimes $f^{-1} = \{(f(a), a) \mid a \in A\}$
is a function from B to A .

Composition function:

If g is a function from B to C
and f " " " " A to B

then $g \circ f = \{(a, c) \mid (\exists b \in B) \text{ with}$
 $f(a) = b \text{ and } g(b) = c\}$

210301) More conditions on functions.

Sketch: U: $g \circ f = \{(a, c) \in A \times C \mid (\exists b \in B)$
 $\left[\underbrace{(a, b) \in f}_{f(a)=b} \wedge \underbrace{(b, c) \in g}_{g(b)=c} \right] \}$ since both are functions!

L: Want: $(\forall a \in A) (\exists! c \in C)$
 $(a, c) \in g \circ f$

equivalently, $\left[\begin{array}{l} \exists c \\ !: \end{array} \right. \forall (a, c), (a, \tilde{c}) \in g \circ f$ and

need $c = \tilde{c}$.

I. Try \exists first:

$\forall a \in A$ have $f(a) = b \in B$

and $g(b) = c \in C$.

so $(a, c) \in g \circ f$ ✓.

For ! if $(a, c), (a, \tilde{c})$ are in $g \circ f$.
then $\exists b, \tilde{b}$ with $(a, b), (a, \tilde{b})$ in f
and $(b, c), (\tilde{b}, \tilde{c})$ in g .

but since f is a function:
have $b = \tilde{b}$
so $(b, c), (b, \tilde{c})$ in g .
so since g is a function:
have $c = \tilde{c}$ similarly. ✓

Ans:

① Not onto since $\text{range}(x^2) = \mathbb{R}_{\geq 0} \neq \mathbb{R}$.
Not one to one since $(-2)^2 = (2)^2 = 4$.

② Not onto since $\text{range}(e^x) = \mathbb{R}_{>0} \neq \mathbb{R}$.

Is one to one since

if $e^x = e^y$

then $\left[\begin{array}{c} \ln(e^x) = \ln(e^y) \\ \parallel \quad \parallel \\ x \quad y \end{array} \right] \checkmark$

Another
proof.

Since e^x is increasing if
 $x < y$ then $e^x < e^y$ so $e^x \neq e^y$.

③ Is onto since $\text{range}(x^3) = \mathbb{R}$.

It's one to one since
if
$$x = \sqrt[3]{x^3} = \sqrt[3]{y^3} = y$$

Def: If f is a function from A to B
which is both one to one and onto
call f bijjective or a bijection.

Ex: ③ above $h(x) = x^3$ is bijjective

From \mathbb{R} to \mathbb{R} .

Claim 4.3.2:

If g a fn. from B to C
and f " " A to B

and $g \circ f$ is an onto fn from A to C

then g is also onto C .

Proof: If c is in C then

$c = (g \circ f)(a)$ for some a in A since $g \circ f$ is onto
and if $b = f(a)$ then $g(b) = g(f(a)) = c$.
q.e.d.

Claim: If $g: B \rightarrow C$ and $f: A \rightarrow B$
and $g \circ f: A \rightarrow C$ is onto.

then $f: A \rightarrow B$ might not be onto.

Proof: Consider the example:

$$A = \{1\}, \quad B = \{2, 3\}, \quad C = \{4\}.$$

$$f(1) = 3, \quad g(2) = g(3) = 4$$

check: $(g \circ f)(1) = 4$ so $g \circ f$ is onto

but $f(1) = 3 \neq 2$

so f is not onto.

210303 §4.4

Def: A function f from A to B is

a bijection or bijjective or

a one to one correspondence

if f is both one to one and onto.
injective surjective.

and if $A = B$ f is called a
permutation.

(§5.1) will use bijections:

Def: If A and B are sets
then $A \approx B$ or A is equivalent
to B if there is a bijection from
 A to B .

Sketch:

$L: A$ reln f from A to B
is a function:

if ① $(\forall a \in A)(\exists b \in B) (a, b) \in f$

and (ii) $(\forall a \in A) (\forall b, \tilde{b} \in B) ((a, b) \in f) \wedge (a, \tilde{b}) \in f \Rightarrow b = \tilde{b}$

or equiv: $(\forall a \in A) (\exists! b \in B) (a, b) \in f$

A function f from A to B

is a bijection:

if (i) onto $(\forall b \in B) (\exists a \in A) (a, b) \in f$

and (ii) 1-1 $(\forall b \in B) (\forall a, \tilde{a} \in A) ((a, b) \in f) \wedge (\tilde{a}, b) \in f \Rightarrow a = \tilde{a}$

or equivalently: $(\forall b \in B) (\exists! a \in A) (a, b) \in f$

I: If f has

fn: $(\forall a)(\exists! b) (a, b) \in f$

$(a, b) \in f$

bij: $(\forall b)(\exists! a)$

then: f^{-1} has:

bij: Since f is a fn so $(a, b) \in f$ so $(b, a) \in f^{-1}$

Fn: since f is a bij so $(a, b) \in f$ so $(b, a) \in f^{-1}$

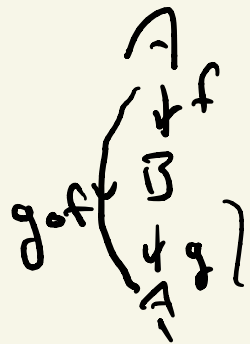
Thm 4.4.4, a:

If f is a fn. from A to B
and g " " " " B to A

then $f = g^{-1}$ iff

$$g \circ f = I_A \quad \text{and} \quad f \circ g = I_B$$

Claim: There are fns:



f from A to B $f \neq f^{-1}$
 and g from B to A B
 with $g \circ f = I_A$ but $f \circ g \neq I_B$.

Proof of Claim: Take $A = \{1\}$, $B = \{2, 3\}$.

and $f(1) = 2$, $g(2) = g(3) = 1$.

Hence $g \circ f(1) = g(2) = 1$ so $g \circ f = I_A$.
 But $f \circ g(3) = f(1) = 2 \neq 3$ so $f \circ g \neq I_B$.
QED.

Proof of Thm:

\Rightarrow Show $g \circ f \subseteq I_A$.

Assume $(a, \tilde{a}) \in g \circ f$ so there is $b \in B$

with $(a, b) \in f$ so $(b, a) \in g$.

and $(b, \tilde{a}) \in g$ so $a = \tilde{a}$ so $g \circ f \subseteq I_A$.

Show $g \circ f \supseteq I_A$.

Assume a is in A

hence there is b in B with

(a, b) in f so (b, a) is in g .

and (a, a) is in $g \circ f$.

Similarly $f \circ g = I_B$.

Remains to check \Leftarrow . (Leave to you).

Sketch

① Compute:

$$\bar{0} \xrightarrow{f} (\bar{0})^2 = \bar{0}$$

$$\bar{1} \xrightarrow{f} (\bar{1})^2 = \bar{1}$$

$$(\bar{2}) \xrightarrow{f} (\bar{2})^2 = \bar{4}$$

$$\bar{3} \xrightarrow{f} (\bar{3})^2 = \bar{9} = \bar{4}$$

Proof of ①: $f(\bar{2}) = \bar{4} = f(\bar{3})$

so f is not one to one.
eqv.

② Compute:

$$0 \xrightarrow{f} 0$$

$$1 \xrightarrow{\quad} 1$$

$$2 \xrightarrow{\quad} \bar{8} = \bar{3}$$

$$3 \xrightarrow{\quad} \bar{27} = \bar{2}$$

$$\bar{1} = \bar{4} \xrightarrow{\quad} \bar{1} = \bar{4}$$

210305

Sketch for claim

③:

U_i

$\mathbb{N} \times \mathbb{N}$

5	9					
4	7	14	28			
3	5	10	20	40		
2	3	6	12	24		
1	1	2	4	8	16	32
\mathbb{N}	1	2	3	4	5	6

f

f is a bijection if every number in \mathbb{N}

appears exactly once in the grid.

L: Is a function!

so need: bijection:

$$- \forall \cancel{x}^n \in \mathbb{N} \quad \exists! (\cancel{a}, \cancel{b}) \in \mathbb{N} \times \mathbb{N}$$

$$\text{with } f(\cancel{(a,b)}) = \cancel{x}^n$$

or

$$\left[\begin{array}{l} \text{exists:} \\ \text{(onto)!} \\ \text{unique!} \\ \text{(1-1)} \end{array} \right. \quad \begin{array}{l} \forall n \in \mathbb{N} \quad \text{there is } (a,b) \\ \text{with } f(a,b) = n \\ \text{If } f(a,b) = f(\tilde{a}, \tilde{b}) \\ \text{then } a = \tilde{a} \text{ and } b = \tilde{b} \end{array}$$

I: (onto): If $n \in \mathbb{N}$

then $n = 2^d \beta$
with $d \geq 0$ and $(\beta \geq 1 \text{ and odd})$

eg $12 = 2^2 \cdot 3$

so $d+1 \geq 1$ is in \mathbb{N}

and $\frac{\beta+1}{2}$ is in \mathbb{N}

and $n = f(d+1, \frac{\beta+1}{2}) = 2^{d+1-1} (2^{\frac{\beta+1}{2}} - 1)$

$$\sqrt{a} \sqrt{b}$$

$$2^{\alpha\beta} \checkmark$$

$$(1-1): \text{ If } f(a, b) = f(\tilde{a}, \tilde{b})$$

$$2^{a-1}(2b-1)$$

$$2^{\tilde{a}-1}(2\tilde{b}-1)$$

For contradiction assume $a \neq \tilde{a}$
 may assume $a > \tilde{a}$

divide both sides by $2^{\tilde{a}-1}$

$$\text{get } 2^{a-\tilde{a}}(2b-1) = (2\tilde{b}-1)$$

so even = odd \times .

Hence ~~if~~ $a = \tilde{a}$ ~~and~~ then
 $2^{a-1}(2b-1) = 2^{a-1}(2\tilde{b}-1)$

so $2b-1 = 2\tilde{b}-1$

so $b = \tilde{b}$ ✓

Notes:

$f: A \rightarrow B$ is a bijection

: $f \circ f$ (~~Thm 4.4.4~~) (Hw 4.4.6)

$\exists g: B \rightarrow A$ with $f \circ g = I_B$ and $g \circ f = I_A$

$; f f$ (Thm 4.4.4)

$\exists g: B \rightarrow A$ with $g = f^{-1}$

$; f f$ (Cor 4.4.3)

$\exists g: B \rightarrow A$ with $g = f^{-1}$
and g is a bijection.

Ans:

$$\Sigma(1, \infty) \xrightarrow{e} (1, \infty) \xrightarrow{b^{-1}} (0, 1)$$

$\not\subset$
 $(0, 1]$ $f = b^{-1} \circ e \circ c$

$$f\left(\frac{1}{2}\right) = b^{-1}\left(e\left(\underbrace{c\left(\frac{1}{2}\right)}_2\right)\right) = \frac{1}{3}$$

$$\underbrace{\underbrace{\quad}_2}_3 = \frac{1}{3}$$

$$f\left(\frac{1}{n}\right) = \frac{1}{n}$$

$$f(1) = \frac{1}{2}$$

$$f(x) = \left\{ \begin{array}{ll} \frac{1}{\frac{1}{x} + 1} = \frac{x}{1+x} & \text{if } \frac{1}{x} \in \mathbb{N} \\ \frac{1}{\frac{1}{x}} = x & \text{else} \end{array} \right.$$

210308 Ch5: Counting and set cardinality:

Def: If A and B are sets then
 $A \approx B$ or A is equivalent to B
if $\exists A \xrightarrow{f} B$ with f a bijective function

Ex: $\overset{\mathbb{N}_3}{\{1, 2, 3\}} \approx \{a, b, c\} \neq \{1, 2, 3\}$

Def: $\mathbb{N}_k = \{1, 2, \dots, k\}$

$\mathbb{N}_0 = \varnothing$

Def: If $A \approx \mathbb{N}_k$ write $\bar{A} = k = \overline{\overline{\{a, b, c\}}}$

and say A has cardinality k .

eg $\overline{\overline{\varnothing}} = 0$, $\overline{\overline{\{a, b, c\}}} = 3$

If $A \approx \mathbb{N}$ write $\overline{A} = \aleph_0$
aleph not

If $A \approx (0, 1)$ write $\overline{A} = c$

→ and say A is denumerable
or countably infinite.

→ and say A has the cardinality
of the continuum.

If $\overline{A} = k$ or \aleph_0 call A countable.

If $\overline{A} = k$ call a finite and
otherwise infinite.

210316 Inequivalence of sets:

Inductive step in proof that-

$P(n) : \forall r < n \quad \exists f: N_n \rightarrow N_r$
which is injective.

Use: [HW problem -
If $x \in N_k$

$\exists g: N_k - \{x\} \rightarrow N_{k-1}$
a bijection.

Assume $P(n)$

Use contradiction to show $P(n+1)$.

Assume for contradiction $\sim P(n+1)$.

so $\exists f: N_{n+1} \longrightarrow N_r$ injective
with $r < n+1$

hence: $f|_{N_n}: N_n \longrightarrow N_r - \{x\}$

if $x = f(n+1)$ since f is injective,

and $f|_{N_n}$ is injective.

and by HW $\exists g: N_r - \{x\} \longrightarrow N_{r-1}$
injective.

so $g \circ f|_{N_n}: N_n \longrightarrow N_{r-1}$ injective.

and $r-1 < n$ so

$\exists g \circ f|_{N_n}: N_n \longrightarrow N_{r-1}$ inj. and $r-1 < n$

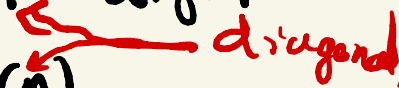
contradicting $P(n)$.

q.e.d.

Thm 5.2.4: $\mathbb{N} \not\approx (0,1)$

Proof Plan: Show $\nexists f: \mathbb{N} \rightarrow (0,1)$ surjective.

Assume $f: \mathbb{N} \rightarrow (0,1)$ and find
 $x \in (0,1)$ with $x \notin \text{Im}(f)$,

Clever diagonalization trick.
consider decimal expansions of each $f(n)$
and write a_n for the n^{th} digit
of the decimal exp of $f(n)$  diagonal.

eg: $f(1) = .\overset{\textcircled{3}}{3}2154 \dots$
 $f(2) = .5\overset{\textcircled{5}}{5}51122 \dots$
 $f(3) = .35\overset{\textcircled{6}}{6}13 \dots$
 $f(4) = .335\overset{\textcircled{5}}{5}1111$
 $f(5) = .9987\overset{\textcircled{5}}{5}443 \dots$

(a_n) $a_1=3$ $b_1=5$
 $a_2=5$ $b_2=3$
 $a_3=6$ $b_3=5$
 $a_4=5$ $b_4=3$
 $a_5=6$ $b_5=3$
 \vdots

choose: $b_n \neq a_n$ eg $b_n = \begin{cases} 5 & \text{if } a_n \neq 5 \\ 3 & \text{if } a_n = 5 \end{cases}$

check $B = .b_1b_2b_3\dots \notin \text{Im}(f)$

eg: $= .53533\dots$

Thm 5.4.3: If A is a set
then $A \neq P(A)$

Proof sketch
w: for A finite:

if $\bar{A} = n$ then $\overline{P(A)} = 2^n \neq n$

eg. $\overline{\emptyset} = 0$ and $\overline{P(\emptyset)} = 2^0 = 1 \neq 0$
 $\overline{\{1\}} = 1$ and $\overline{P(\{1\})} = 2^1 = 2 \neq 1$

L: for any $f: A \rightarrow P(A)$
show f is not onto

and so not bijective.

Trick/Idea: Given f build.

$$B = \{a \in A \mid a \notin f(a)\}$$

eg: $f: \{1, 2, 3\} \rightarrow \mathcal{P}(\{1, 2, 3\})$,
 $f(1) = \{1, 2\}$, $f(2) = \{3\}$, $f(3) = \{1, 2, 3\}$.

$$B = \{2\}$$

Key: $B \notin \text{Im}(f)$

Note: $(0,1) \neq P((0,1)) \neq P(P((0,1)))$
 $\neq P(P(P((0,1)))) \dots$

210312

First order Predicate Logic:

Propositions:

Truth Tables

inside back cover

$\wedge, \vee, \sim, \Rightarrow, \Leftrightarrow$

universe

\exists

\forall

$\exists!$

Sets:

Venn Diagrams

$\cap, \cup, \underbrace{A^c}_{\text{in a universe}}, \in, \bigcap_{\text{family}}, \bigcup_{\text{family}}, \times, -$

$\emptyset, \mathcal{P}(A), \overline{\overline{A}}, \approx$

Relations: $R \subseteq A \times B$

$x R y$ or $(x, y) \in R$

R^{-1} , $R \circ S$, $\text{dom}(R)$, $\text{rang}(R)$.

Equivalence relations:
3 properties

\bar{x} , A/R (set of sets).

eg $\mathbb{Z}_n = \mathbb{Z}/R$ for the right R .

Functions: $\forall x \exists! y$ with $(x, y) \in f$ with $f(x) = y$

bijective functions: $\forall y \exists! x$ with $(x, y) \in f$.
2 properties

Example Problems:

Find $\gcd(182, 21)$

Ans: Two approaches:

① Euclid's Alg:

$$182 = 8 \cdot 21 + 14$$

$$21 = 1 \cdot 14 + \boxed{7} = \text{gcd}$$

$$14 = 2 \cdot 7$$

② Factor both: $2 \cdot \boxed{7} \cdot 13 = 182$
 $3 \cdot \boxed{7} = 21$
 $\quad \quad \quad = \text{gcd.}$

Problem:

(a) Prove that $\underbrace{S = \{(x, y) \in \mathbb{R}^2 \mid x - y \in \mathbb{Q}\}}_{\text{relation on } \mathbb{R}}$

is an equivalence relation.

(b) Find $x, y, z \in \mathbb{R}$ with $\overline{x} = \overline{y} \neq \overline{z}$

(c) Show that $\overline{\pi}$ is denumerable.

$$\textcircled{b} \quad \overline{x} = \{y \mid (x, y) \in S\}$$

eg $\overline{x} = \overline{y}$ if $x = y$

eg $\overline{0} = \overline{\frac{1}{2}}$

Proof:

\subseteq if $a \in \overline{0}$ then $0 - a \in \mathbb{Q}$ so $\frac{1}{2} - a \in \mathbb{Q}$ so $a \in \overline{\frac{1}{2}}$.


$(0, a) \in S$ so

$\frac{1}{2} - a \in \mathbb{Q}$

so $(\frac{1}{2}, a) \in S$

so $a \in \overline{\frac{1}{2}}$.

\supseteq similarly,

and  $\bar{0} \neq \bar{\pi}$

since: $\pi \in \bar{\pi}$ since $(\pi, \pi) \in S$
since S is reflexive
(or $\pi - \pi = 0 \in \mathbb{Q}$).

but $\pi \notin \bar{0}$ since

for contradiction assume $\pi \in \bar{0}$

so $(0, \pi) \in S$ so $0 - \pi \in \mathbb{Q}$.

but $-\pi \notin \mathbb{Q}$. a contradiction.

qed.

