THE NUMBER OF POLYTOPES, CONFIGURATIONS AND REAL MATROIDS

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Abstract. We show that the number of combinatorially distinct labelled d-polytopes on n vertices is at most $(n/d)^{d^2n(1+o(1))}$, as $n/d \to \infty$. A similar tool is a theorem of Milnor and Thom from real algebraic geometry real oriented and unoriented matroids with n elements of rank d. Our mair bounds. We also obtain sharp upper and lower bounds for the numbers of Goodman and Pollack. This bound improves considerably the previous known bound for the number of simplicial polytopes has previously been proved by

easy; there is only one polygon with n (unlabelled) vertices and there are solve the general problem. Write $\beta = n - d$. The cases $d \le 2$ or $\beta \le 2$ are quite c(n, d) and $c_s(n, d)$ for small values of n and d, however, does not, of course and Altshuler and Steinberg [AS1, AS2] determined c(8, 4). Determining $c_s(11,3)$. More recently, Grünbaum and Sreedharan [GS] determined $c_s(8,4)$ to extend Brückner's work for n = 11, 12, but both his enumeration and attempts were given by Brückner [Br] and Steinitz [Ste] (see also [Gr. pp. c(n,3) or $c_s(n,3)$ despite a lot of effort. Detailed historical surveys of these 91-92]), it seems extremely difficult actually to determine this number even much effort and frustration of nineteenth-century geometers. Although it estimating these two functions (especially for 3-polytopes) was the subject of simplicial d-polytopes on n labelled vertices. The problem of determining on of) d-polytopes on n labelled vertices and let $c_s(n, d)$ denote the number of (see [Gr. pp. 98-101]). $[d^2/4]$ d-polytopes on d+2 (unlabelled) vertices, [d/2] of which are simplicial [Gra]. Hermes [He] determined c(n, 3) for $n \le 8$ and Grace [Gra] determined Brückner's extensive attempts to correct it were incomplete, as shown by Grace 288-290]). Brückner [Br] determined $c_i(n, 3)$ for $n \le 10$. Hermes [He] tried for relatively small n and d. Both Cayley and Kirkman failed to determine the real field that the problem of computing c(n, d) is solvable (cf. [Gr. pp follows from Tarski's Theorem on the decidability of first order sentences in §1. Introduction. Let c(n, d) denote the number of (combinatorial types

as d tends to infinity. An explicit formula for c(d+3,d) was given later by formula for $c_1(d+3,d)$ and determined the asymptotic behaviour of c(d+3,d)Using a Gale Diagram, Perles (cf. [Gr. pp. 112-114]) found an explicit

precisely by Tutte [Tu] and by Richmond and Wormald [RW], (see also [Gr nation of $c_i(n, d)$ or c(n, d) for $d \ge 4$ and $n \ge d + 4$ is a problem of an entirely pp 289-290]). However, as mentioned in [Gr. p 290], it seems that the determi-The asymptotic behaviour of c(n, 3) and $c_i(n, 3)$ was determined almost

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 $c_s(n,d)$ was $n^{cn^{d/2}}$. This follows easily from the upper bound theorem different order of magnitude. Until recently, the best general upper bound for is a triangulated (d-1) sphere, but the converse is false when $d, \beta \ge 4$). (d-1)-spheres. (Recall that the boundary complex of a simplicial d-polytope [Kl, M, St], and the argument applies also to bound the number of triangulated

the number of simple configurations of n points in R^d is less than $n^{d(d+1)n}$ points is the same as that of their images. By a clever use of a theorem of if there is a bijection $\phi: A \to B$ such that the orientation of each d+1 (ordered) in general position in \mathbb{R}^d . Two such configurations A and B are isomorphic Milnor [Mi] from real algebraic geometry Goodman and Pollack showed that [GP2]. A simple configuration of n points in R^d is an ordered n-tuple of points show that this number is at least This is close to the truth at least for fixed d and large n since it is easy to A major development was very recently achieved by Goodman and Pollack

$$n^{d^2n(1+O(\log d/\log n))}$$
.

considerably the best previously known bound. Moreover, their result gives immediately that $c_s(n, d) \le n^{d(d+1)n}$, improving

also slightly improve the bound of [GP2] and show, in particular, that for and hence to bound the number of arbitrary d-polytopes on n-vertices. We bound the number of simple and nonsimple configurations of n points in R^d , in \mathbb{R}^d both have the form $n^{d^2\pi(1+o(1))}$ as $n\to\infty$. For polytopes we obtain fixed $d \ge 2$ the numbers of simple or of nonsimple configurations of n points In this paper we apply another (similar) theorem of Milnor and Thom to

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$$\left(\frac{n-d}{d}\right)^{nd/4} \leq c_s(n,d) \leq c(n,d) \leq (n/d)^{d^2n(1+O(\log\log(n/d)/\log(n/d)))},$$

of simplicial polytopes. on n vertices is at least 22nd Very recently, Kalai [K] showed that the total number of triangulated spheres and show that the total number of polytopes on *n* vertices is at most $2^{n^3+O(n^2)}$ Thus, very few of these are boundary complexes

and $n \to \infty$, these numbers have the form $n^{O(d^2n)}$. The total number of complex matroids on n elements is bounded by $2^{O(n^2)}$, a very small part of the total number of real and complex matroids with n elements of rank d. For fixed dnumber of matroids on n points which is at least $2^{\Omega(2^n/n^{3/2})}$, as shown by Knuth Our methods also enable us to obtain sharp bounds on the asymptotic

of n points in \mathbb{R}^d . In the final Section 5 we prove our bounds for the number matroids on n points, and in Section 4 we consider the number of configurations polynomials. In Section 3 we deal with the number of real and complex to obtain a general bound on the number of sign patterns of a sequence of of d-polytopes Our paper is organized as follows: in Section 2 we apply Milnor's Theorem

 $(j=1,\ldots,m)$ be real polynomials. For a point $c=(c_1,c_1,\ldots,c_n)\in R^n$ the sign pattern of the P_i at c is the m-tuple $(\epsilon_1, \epsilon_2, \dots, \epsilon_m) \in \{-1, 0, 1\}^m$, where \$2. The Number of Sign Patterns. Let $P_i = P_i(x_1, x_2, ...)$

 $\varepsilon_j = \operatorname{sign} P_j(c_1, c_2, \ldots, c_n)$. The total number of sign patterns as c ranges over all points of R^n , denoted by $s(P_1, \ldots, P_m)$, is clearly at most 3^m . Using a theorem of Milnor [Mi] (see also Thom [Th]) from real algebraic geometry, we bound this number by a function of n and the degrees of the P_j . All our upper bounds in the paper are derived from this bound (and its analogue for complex polynomials).

We first state Milnor's theorem.

THEOREM 2.1 (Milnor [Mi, Theorem 2]). Let V be a variety in \mathbb{R}^{1} , defined by the polynomial equations

$$f_i(x_1, x_2, ..., x_l) = 0, \qquad (i = 1, ..., h).$$

If each polynomial f_i has degree $\leq k$, then the sum of the Betti numbers of V is at most $k(2k-1)^{l-1}$. In particular, the number of connected components of V is at most $k(2k-1)^{l-1}$.

Using this theorem we prove

THEOREM 2.2. Let $P_1 = P_1(x_1, \ldots, x_n), \ldots, P_m = P_m(x_1, \ldots, x_n)$ be real polynomials. Let $d_i = \deg P_j(\geqslant 1)$ be the degree of P_j , $1 \le j \le m$. Put $J = \{1, 2, \ldots, m\}$ and let

$$J=J_1\cup J_2\cup\ldots\cup J_k$$

be a partition of J into h pairwise disjoint parts. Define

$$k=4 \max_{1 \le i \le h} \left(\sum_{j \in J_i} d_j \right).$$

Then the number of sign patterns of the P_j satisfies

$$s(P_1, ..., P_m) \le k(2k-1)^{n+h-1}$$

Remark 2.3. By taking the trivial partition of J into one part, we conclude that if $r = \sum_{j=1}^{m} d_j$ then $s(P_1, \ldots, P_m) \le 4r(8r-1)^n$. By using another theorem of Milnor [Mi, Theorem 3], we can show that in fact

$$s(P_1, \dots, P_m) \le (2+2r)(1+2r)^{n-1}.$$
 (2.1)

For our applications, however, Theorem 2.2 will usually give asymptotically better bounds. We omit the detailed proof of (2.1).

Proof of Theorem 2.2. Let $C \subseteq R^n$ be a finite set of points that represents all the sign patterns of the P_i . (Clearly there is such a C satisfying $|C| \le 3^n$). For $c = (c_1, c_2, \ldots, c_n) \in C$ and $1 \le i \le m$ we denote $P_j(c_1, \ldots, c_n)$ by $P_j(c)$. Define $\varepsilon > 0$ by

 $\varepsilon = \frac{1}{2} \min \{ |P_j(c)| : c \in C, 1 \le j \le m \text{ and } P_j(c) \ne 0 \}.$

Let $\delta > 0$ satisfy $\delta < \varepsilon^{4/3}$, $1 \le i \le h$. Define h polynomials f_1, f_2, \ldots, f_h with variables $x_1, x_2, \ldots, x_h, y_1, y_2, \ldots, y_h$ by

$$f_i(x_1, \dots, x_n, y_1, \dots, y_n)$$

$$= -y_i^n - \delta + \prod_{i \in I_i} (P_i(x_1, \dots, x_n) - \epsilon)^n (P_i(x_1, \dots, x_n) + \epsilon)^n.$$

The degree of f_i is clearly $4\sum_{j\in J_i}d_j \le k$. Also, if $c=(c_1,\ldots,c_n)\in C$ then

$$\prod_{i \in J_i} (P_j(c) - \varepsilon)^2 (P_j(c) + \varepsilon)^2 > \delta,$$

and hence there exist real values b_1, \ldots, b_h such that

$$f_i(c_1,\ldots,c_n,b_1,\ldots,b_h)=0, \qquad (1 \le i \le h).$$

For $c \in C$ denote by \bar{c} a real vector $(c_1, \ldots, c_n, b_1, \ldots, b_h)$ that satisfies the last system. Let V be the variety in \mathbb{R}^{n+h} defined by

$$f_i(x_1, ..., x_n, y_1, ..., y_h) = 0$$
 $(1 \le i \le h)$

By definition every vector $\bar{c}(c \in C)$ is a point of V. We now claim that if $c_1, c_2 \in C$ represent distinct sign patterns of the P_j , then \bar{c}_1, \bar{c}_2 are not in the same connected component of V. Indeed, if $c_1, c_2 \in C$ represent distinct sign patterns, there exists some $1 \le j \le m$, such that sign $P_j(c_1) \ne \text{sign } P_j(c_2)$. If $j \in J_i$, this implies, by continuity and the choice of ε , that any path in R^{n+h} joining \bar{c}_1 to \bar{c}_2 contains a point $(x_1, \ldots, x_n, y_1, \ldots, y_h)$ such that $P_j(x_1, \ldots, x_n) = \varepsilon$, or $P_j(x_1, \ldots, x_n) = -\varepsilon$, i.e., a point where

$$f_i(x_1,...,y_h) = -y_i^2 - \delta \le -\delta \le 0.$$

This point is thus not in V and our claim follows. Since C represents all the sign patterns of the P_j , we conclude that the number of sign patterns is at most the number of connected components of V which is, by Milnor's Theorem (Theorem 2.1), at most $k(2k-1)^{n+h-1}$. This completes the proof of the theorem.

For our applications we will also be interested in the number of sign patterns of complex polynomials. If $Q_i(z_1,\ldots,z_n)$, $1 \le i \le m$, are complex polynomials and $b=(b_1,\ldots,b_n)\in C^n$, the sign pattern of the Q_i at b is the m-tuple $(\epsilon_1,\ldots,\epsilon_n)\in\{0,1\}^m$, where $\epsilon_j=\text{sign}\,|Q_j(b)|$. By applying Theorem 2.2 to the real polynomials $P_j=(\text{Re }Q_j)^2+(\text{Im }Q_j)^2$, $1 \le j \le m$ in the 2n real variables Re z_i and Im z_i , $1 \le j \le n$, we can bound the number of sign patterns of the Q_i in terms of their degrees. Moreover, since in this case $P_j \ge 0$, we can slightly improve the estimate by defining here $f_i=-y_i^2-\delta+\prod(P_j-\epsilon)^2$. This gives the following theorem, whose detailed proof is omitted.

THEOREM 2.4. Let $Q_1(z_1,\ldots,z_n),\ldots,Q_m(z_1,\ldots,z_n)$ be complex polynomials. Put $d_j=\deg Q_j(\geqslant 1),\ J=\{1,2,\ldots,m\}$ and let $J=J_1\cup J_2\cup\ldots\cup J_h$ be a partition of J into h pairwise disjoint parts. Define

$$k=4 \max_{1 \le i \le h} \left(\sum_{j \in J_i} d_j\right).$$

Then the number of sign patterns of the Q is at most

$$k(2k-1)^{2n+b-1}$$

§3. The Number of Rational, Real and Complex Matroids. There are several known asymptotic estimates for the number of nonisomorphic matroids of several kinds on n points. See [We, pp. 305-308] for bounds on the number of all matroids on n points, the number of transversal matroids on n points, and the number of matroids on n points which are representable over a finite

matroids on n points is at least field with q elements. Knuth [Kn] showed that the number of labelled (simple)

$$2^{(r_n r_{2j})/(2n)}$$
.

respectively. Clearly are representable over the rationals, the reals and the complex numbers r(n, d, C), which are the numbers of matroids on n points with rank d, which Here we obtain sharp bounds on the numbers r(n d, Q), r(n, d, R) and

$$r(n, d, Q) \le r(n, d, R) \le r(n, d, C).$$

Here we show that

$$n^{(d-1)^2n - O(d^2n(\log d + \log\log n)/\log n)} \le r(n, d, Q) \le r(n, d, R)$$

$$\le n^{(d-1)dn + O(nd \log\log n/\log n)}, \tag{3.1}$$

that

$$r(n,d,C) \le n^{2(d-1)dn+O(nd\log\log n'\log n)},\tag{3.2}$$

and that for every $d \le n$

$$r(n, d, R) \le r(n, d, C) \le 2^{O(n^3)}$$
. (3.3)

total degree in the dn variables x_y . Divide the polynomials into $[n/\log n]$ groups, each of matroids is at most the number of sign patterns of the $\binom{n}{d}$ degree d polynomials number of oriented real matroids—see Section 4). Thus the total number of the sign pattern determines more, and we get here an upper bound on the x_{ij} determines the matroid represented by the given values of the x_{ij} . (In fact, Hence the sign pattern of these $\binom{a}{d}$ polynomials of degree d in the dn variables is non-zero, if, and only if, the corresponding set is a base of the matroid $i \in \{i_1, \dots, i_d\}, j = 1, \dots, d$, where $1 \le i_1 < i_2 < \dots < i_d \le n$. Such a determinant our matroid, and consider the set of all $\binom{n}{d}$ d by d determinants $\det(x_{ij})$, is independent, if, and only if, the corresponding vectors are linearly independent. Let $(x_{11}, \ldots, x_{1d}), \ldots, (x_{n1}, \ldots, x_{nd}) \in \mathbb{R}^d$ be the vectors representing i.e., for each point of the matroid we have a vector in R^d and a set of points It is easy to check that every real matroid of rank d is representable in R^d We first prove the upper bounds. We begin by considering real matroids.

$$\sim \binom{n}{d} d/[n/\log n] \leq (1+o(1)) \frac{n^{d-1}\log n}{(d-1)!}$$

and apply Theorem 2.2 with

$$h = [n/\log n], \quad k = (4 + o(1)) \frac{n^{d-1} \log n}{(d-1)!}$$

polynomials is at most $\binom{n}{d} d \le 2^n n$ (or Theorem 2.2 with h=1) by noticing that the sum of degrees of our follows similarly from Theorem 2.4. Inequality (3.3) follows from Remark 2.3 (and dn variables) to get the upper bound in inequality (3.1). Inequality (3.2)

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The only remaining part is the lower bound in equality (3.1). It suffices

$$r(n, d, Q) \ge \left(\frac{[n/\log n]}{(d-1)^2} \right)^{n-[n/\log n]}$$

space of dimension d-1, with all coordinates having rational values. Assume, $p(N) \in Q^{d-1}$ (since it is a solution of a linear system of equations with rational choices of the N_i yield different points p(N). Moreover, as is easily checked, of our points each H_i is uniquely defined and p(N) is a point. Also, different and so on. Let H_i be the hyperplane of R^{d-1} containing the points in N_i and first (d-1) points of N, the second consisting of the next (d-1) points of N $... \cup N_{d-1}$ be a partition of N into d-1 equal parts, the first consisting of the $|N| = (d-1)^2$, we define a point $p(N) \in Q^{d-1}$ as follows. Let $N = N_1 \cup N_2 \cup N_3 \cup N_4 \cup N_4$ Let M be a labelled set of $[n/\log n]$ points in Q^{d-1} , i.e., in the real Euclidean each being one of the p(N). There are coefficients). We now add to M another labelled set of $n - \lfloor n/\log n \rfloor$ points let p(N) be the intersection point of these hyperplanes. By the generic position further, that the set M is in a generic position. For every subset $N \subseteq M$,

$$\left(\begin{bmatrix} n/\log n \end{bmatrix}\right)^{n-[n/\log n]} = h(n, d)$$

all our h(n, d) labelled sets of points supply distinct matroids, and the lower n points with rank d, in which $\{i_1,\ldots,i_d\}$ is independent if x_{i_1},\ldots,x_{i_d} span $y_i = (x_1, \dots, x_{i(d-1)}, 1)$. The y_i form a representation of a rational matroid on possibilities for this construction, each supplying a set of n labelled points in bound of inequality (3.1) follows. R^{d-1} , i.e., are not contained in a hyperplane in R^{d-1} . It is easy to check that Q^{d-1} . If $x_i = (x_{i1}, \dots, x_{i(d-1)})$ are the coordinates of the *i*-th point, put

 $0 \le k, j \le d$, where $P_{i_1} = (x_{i_1}, \dots, x_{i_k d})$ and $x_{i_k o} = 1$ for $0 \le k \le d$. The order type of a configuration C of points is sometimes known as the oriented matroid of n labelled points in \mathbb{R}^d . In [GP1] Goodman and Pollack showed that the number of distinct order types of configurations of n labelled points in structure determined by C. (See [GP1] for more details.) Let t(n, d) denote is simple if $w(S) \neq 0$ for every such S. Notice that w(S) is just sign det $(x_{i,j})$. w(S) = -1 if $P_{i_1} \dots P_{i_d} < 0$, and w(S) = 0 if $P_{i_1} \dots P_{i_d} = 0$. The configuration $S = \{i_0, i_1, \dots, i_d\}$ with $1 \le i_c < i_1 < \dots < i_d \le n$, w(S) = +1 if $P_{i_1} \dots P_{i_d} > 0$, w from the set of all (d+1)-subsets of $\{1, 2, ..., n\}$ to $\{0, \pm 1\}$, where for type of a configuration C of n labelled points P_1, P_2, \ldots, P_n in \mathbb{R}^d is a function orientation, written $P_o \dots P_d > 0$, if det $(x_{ij})_{0 \le i,j \le d} > 0$ where $x_{io} = 1$ for each i. points in R^d , with $P_i = (x_{i1}, \dots, x_{id})$ for each i, we say they have positive §4. The Number of Configurations. If (P_0, P_1, \ldots, P_d) is a sequence of d+1of Milnor (mentioned in Remark 2.3 above) to prove that $t_s(n,d) \le n^{n^\alpha}$. Very recently [GP2], they found a clever way of using a theorem The conditions $P_o \dots P_d < 0$ and $P_c \dots P_d = 0$ are defined similarly. The *order* R^d , and let $t_i(nd)$ denote the number of order types of simple configurations

$$I_{s}(n,d) \leq n^{d(d-1)n}.$$

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This bound has several interesting applications (see [GP2], [AFR]). Goodman and Pollack also showed that

$$I_s(n,d) \ge n^{d^2n(1+O(\log d)\log n)}. \tag{4.2}$$

Here we apply Theorem 2.2 to show that the total number of order types t(n, d) is not much bigger than $t_s(n, d)$. In fact we also slightly improve (4.1) and prove the following theorem, which supplies very sharp estimates for the asymptotic behaviour of both functions t(n, d) and $t_s(n, d)$, at least for n much greater than d.

THEOREM 4.1.

$$(n/d)^{d^2n + O(d^2n \frac{\log d}{\log n})} \le t_s(n,d) \le t(n,d)$$

$$\le (n/d)^{d^2n(1+O(\frac{1}{\log (n/d)} + \frac{\log \log (n/d)}{d \log (n/d)}))},$$

Proof. The lower bound is just inequality (4.2). (Notice that $d^{d^2n} = n^{d^2n \log d/\log n}$). To get the upper bound, notice that r(nd) is just the number of sign patterns of $\binom{d^2n}{d^2n}$ polynomials of degree d in the dn real variables (x_{i1}, \ldots, x_{id}) , which are the coordinates of the i-th point. The polynomials are just all the determinants $\det(x_{i,j})$, $0 \le k$, $j \le d$, where $x_{i,ko} = 1$ for all k and $1 \le i_0 < i_1 \ldots < i_d \le n$. Split these polynomials into $h = \lfloor n/\log (n/d) \rfloor$ classes, each of total degree

$$\sim \binom{n}{d+1} d \frac{\log (n/d)}{n} \leq \frac{n^d}{d!} \log (n/d).$$

Apply Theorem 2.2 with

$$h = [n/\log (n/d)], \qquad k = 4 \frac{n^d \log (n/d)}{d!}$$

(and dn variables) to conclude that

$$t(n,d) \le (2k)^{\frac{d}{dn} + \frac{n}{\log n + d}} \le \left(8\left(\frac{en}{d}\right)^d \log\left(n/d\right)\right)^{\frac{d}{dn} + \frac{n}{\log n + d}}$$
$$= (n/d)^{\frac{d}{dn}} (1 + O(\frac{1}{\log\left(n/d\right)} + \frac{\log\log\left(n/d\right)}{d\log\left(n/d\right)}))$$

Theorem 4.1 implies that if d and n vary and $\log d/\log n \to 0$, then both $t_i(n, d)$ and t(n, d) have the form

$$(n'd)^{d^2n(1-o(1))} = n^d 2^{n(1+o(1))}. \tag{4.3}$$

In particular we obtain the following.

Corollary 4.2. For fixed d > 2, as $n \to \infty$

$$I_s(n,d) = (n/d)^{d^2n(1+o(1))} = n^{d^2n(1+o(1))}$$

 $t(n,d) = (n/d)^{d(n)(1-\alpha)!} = n^{d(n)(1-\alpha)!}$

and

With a somewhat more careful computation one can extend the range in which $t_s(n, d)$ and t(n, d) have the form (4.3). The most important cases, however, seem to be these covered by Corollary 4.2.

Remark 4.3. By a similar application of Theorem 2.2 one can easily show that for every n and d,

$$t_s(n, d) \le t(n, d) \le 2^{n^3 + O(n^2)}$$

We omit the details.

Remark 4.4. A linear space P on a set $N = \{1, 2, ..., n\}$ of n points is a family L of subsets (called lines) $l_1, ..., l_r$ of N, such that every two points belong to a unique line. If P can be realized by embedding the points in the plane R^2 where L is the set of all maximal collinear subsets of N, P is called a representable linear space. By Corollary 4.2 (with d=2) the number of distinct representable linear spaces on n labelled points is at most $n^{(4+c(1))n}$. On the other hand, it is easy to see that the number of distinct linear spaces on n labelled points is much bigger—it is at least $2^{(n^2/6)+O(n)}$. Indeed, take a fixed Steiner triple system on $\ge n-3$ of our points and let

$$B_1, B_2, \ldots, B_m \qquad (m = n^2/6 + O(n))$$

be the set of its blocks. For every subset F of the set of these blocks, we define a linear space on N whose lines are all the blocks in F together with all pairs of points $\{i,j\}$ that do not lie in any common block of F. This supplies $2^{(n^2/6)+O(n)}$ distinct linear spaces on N.

There are obvious generalizations of this remark to higher dimensions.

§5. The Number of Convex Polytopes. Let c(n,d) denote the number (up to combinatorial isomorphism) of d-polytopes on n labelled vertices and let $c_s(n,d)$ be the number of simplicial d-polytopes on n labelled vertices. The problem of determining or estimating c(n,d) and $c_s(n,d)$ has a long history, part of which, including some previous results, is outlined in Section 1. Very recently, Goodman and Pollack [GP2] used their bound for $t_s(n,d)$ (see inequality (4.1) above) to show that

$$c_s(n,d) \leq t_s(n,d) \leq n^{d(d+1)n}$$

This follows immediately from the fact that the two vertex sets of two inequivalent simplicial polytopes with vertices in general position in R^d form distinct simple configurations. Indeed, one can easily check (see, e.g., [GP1]) that the order type of a configuration that spans R^d determines which sets of its points lie on supporting hyperplanes of its convex hull. This also holds for non-simple configurations. Hence, the order type of a configuration on a set $N = \{1, 2, ..., n\}$ of n points in R^d which is the set of vertices of a convex polytope P determines its facets and thus its complete combinatorial type. This implies that $c_s(n, d) \le c(n, d) \le t(n, d)$, and by Theorem 4.1 and the remarks following it we obtain:

THEOREM 5.1.
$$c_r(n,d) \le c(n,d) \le (n/d)^{\frac{d^2n}{(1+O(\frac{1}{\log(n/d)} + \frac{\log\log(n/d)}{d\log(n/d)}))}}$$

In particular, if
$$n/d \to \infty$$
 then
$$c_s(n,d) \le c(n,d) \le (n/d)^{d(n/1+O(1))}.$$
(5.1)

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Furthermore, for every d and

$$c(n,d) \leq 2^{n^2+O(n^2)},$$

and hence the total number of polytopes on n points is at most $2^{n^2+O(n^2)}$

and c(n, d) by (5.1) is not so far from the truth. Indeed, one can show that for $n \ge 2d$ As mentioned in [GP2], one can show that the estimate given for $c_r(n, d)$

$$c_s(n,d) \geqslant \left(\frac{n-d}{d}\right)^{nd/4} \tag{5.2}$$

magnitude as the total number of d-polytopes on n vertices; quite a surprising polytopes with n points in R^d is $\ge n^{c_d n}$, where $\lim_{d\to\infty} c_d = 1/2$. By (5.1) this one neighbourly d-polytope on n points). fact (especially in view of Motzkin's old conjecture [Mo] that there is only shows that for fixed $d(\ge 4)$ this number is of roughly the same order of Shemer proved that even the number of (unlabelled) distinct neighbourly possibilities, each one "close" to a facet of P. This implies (5.2). In [Sh] P has $\ge ((n-d)/d)^{d/2}$ facets. Put the last n/2 labelled points, in all To see this, take a cyclic polytope P on the first n/2 points (see [Gr]). Then

n-d=o(n). In fact, by being more careful we can prove that, for fixed $\beta>0$ Theorem 4.1 and Theorem 5.1 can be somewhat improved for the case We conclude our paper by noting that, as observed by G. Kalai, both

$$c_{\varepsilon}(d+\beta,d) \leq c(d+\beta,d) \leq t(d+\beta,d) \leq n^{\beta(\beta-1)(d(1+\alpha(1)))}$$

We omit the details.

discussions. I also thank G. Kalai for many extremely helpful suggestions. Acknowledgment. I thank J. Goodman and R. Pollack for stimulating

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Received on the 17th of July, 1985 geometry: Other problems of combinatorial RELATED GEOMETRIC TOPICS; CONVEX SETS AND

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