

MAT 145, Spring 2020
Solutions to homework 7

1. (20 points) Prove that a graph with n nodes and m edges ($m < n$) has at least $n - m$ connected components.

Solution 1: We prove it by induction in m . If $m = 0$ then we have n vertices and no edges, so n connected components. Suppose that we proved this for a graph with m edges, consider a graph G with $(m + 1)$ edges. Let G' be a subgraph of G with one edge e removed. By the assumption of induction, G' has k connected components and $k \geq n - m$. The edge e either connects two components of G' and then G has $k - 1$ components, or it connects two vertices in the same connected components of G' and then G has k components. In any case, G has at least $k - 1 \geq n - m - 1 = n - (m + 1)$ components.

Solution 2: If G is a connected graph then it has at least $n - 1$ edges. Indeed, a spanning tree in G has exactly $n - 1$ edges.

Now suppose that G has k connected components with n_1, \dots, n_k vertices and m_1, \dots, m_k edges. Then by the above $m_i \geq n_i - 1$, so

$$m = m_1 + \dots + m_k \geq (n_1 - 1) + \dots + (n_k - 1) = (n_1 + \dots + n_k) - k = n - k.$$

Therefore $m \geq n - k$ and $k \geq n - m$.

2. (20 points) Prove that if a tree has a node of degree d , then it has at least d leaves.

Solution: since we are in a tree, any two vertices are connected by a unique path. Suppose that v has d neighbors v_1, \dots, v_d . Any vertex u of distance n from v has exactly one neighbor which is distance $(n - 1)$ from v (along the path from v to u) and all other neighbors have distance $(n + 1)$ from v .

Consider d paths in the tree which are constructed as follows. We start from v , then go to v_i , then go to an arbitrary neighbor of v_i which is distance 2 from v , then go to an arbitrary neighbor of that vertex which is distance 3 from v and so on. This process will eventually stop in a leaf, and all these leaves are different for different i . Therefore, there are at least d different leaves.

3. (20 points) A rooted binary tree has height n , that is, the distance from every vertex to the root is at most n . What is the maximal number of vertices in this tree?

Solution: We have 1 root (distance 0 from the root), 2 vertices of distance 1 from the root, 2^2 vertices of distance 2 from the root, ... 2^n

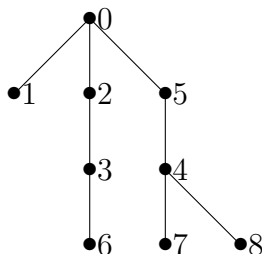
vertices of distance n from the root. In total there are $1 + 2 + \dots + 2^n = 2^{n+1} - 1$ vertices.

4. (20 points) A tree with 9 vertices (labeled from 0 to 8) has second row in Prüfer code 03204450. Reconstruct the first row in Prüfer code and the tree.

Solution: By Prüfer's algorithm we reconstruct the first row in the code:

$$\begin{array}{cccccccc} 1 & 6 & 3 & 2 & 7 & 8 & 4 & 5 \\ 0 & 3 & 2 & 0 & 4 & 4 & 5 & 0 \end{array}$$

and the tree



5. (20 points) Prove that the vertices in a tree can be colored black and white such that no two vertices of the same color are connected by an edge.

Solution 1: By a theorem from lecture, a graph can be colored in black and white if and only if it does not have a cycle of odd length. Since there are no cycles in a tree, it can be colored.

Solution 2: We prove by induction by the number of vertices n . For $n = 1$ it is clear. Suppose that we proved it for any tree with n vertices, consider a tree T with $(n + 1)$ vertices. Since T is a tree, it has a leaf v . Then $T - v$ is a tree with n vertices, and can be colored by the assumption of induction. If the neighbor of v in $T - v$ is colored white, we color v black, if the neighbor of v in $T - v$ is colored black, we color v white. Since v is a leaf, it has no other neighbors and it is a valid coloring.