## CSE 6321 - Problem Set 3 Due lecture on April 18th

Collaboration is permitted; looking for solutions from external sources (books, the web, material from previous years, etc.) is prohibited.

- 1. Show that if every NP-hard language is also PSPACE-hard, then PSPACE = NP.
- 2. Show that TQBF restricted to formulas where the part following the quantifiers is in conjunctive normal form is still PSPACE-complete.
- 3. Read the definition of MIN-FORMULA from Problem set 2.
  - (a) Show that  $MIN FORMULA \in PSPACE$ .
  - (b) Explain why this argument fails to show that  $MIN-FORMULA \in coNP$ : If  $\phi \notin MIN FORMULA$ , then  $\phi$  has a smaller equivalent formula. A NTM can verify that  $\phi \in \overline{MIN FORMULA}$  by guessing that formula.
- 4. An undirected graph is *bipartite* if its nodes may be divided into two sets so that all edges go from a node in one set to a node in the other set. Show that a graph is bipartite if and only if it doesn't contain a cycle that has an odd number of nodes. Let  $BIPARTITE = \{\langle G \rangle : G \text{ is a bipartite graph}\}$ . Show that  $BIPARTITE \in NL$ .
- 5. Recall that a directed graph is *strongly connected* if every two nodes are connected by a directed path in each direction. Let

 $STRONGLY-CONNECTED = \{\langle G \rangle : G \text{ is a strongly connected graph}\}.$ Show that STRONGLY - CONNECTED is NL-complete.