

## CSE 6321 - Problem Set 2

Due beginning of lecture on February 16th

Problem numbers are from the third edition of Sipser's book. If unsure about which problem to solve, ask. Collaboration is permitted; looking for solutions from external sources (books, the web, material from previous years, etc.) is prohibited. Printed version is preferred, otherwise please make sure your handwriting is readable.

1. Let  $C \subseteq \Sigma^*$  be a language. Prove that  $C$  is Turing-recognizable iff a decidable language  $D$  exists such that  $C = \{x \in \Sigma^* : (\exists y \in \Sigma^*) \langle x, y \rangle \in D\}$ .
2. Let  $T = \{\langle M \rangle : M \text{ is a T.M. that accepts } w \text{ reversed whenever it accepts } w.\}$ . Determine whether  $T$  is decidable, undecidable but recognizable or unrecognizable. Prove your answer.
3. Let

$$M = \{\langle a, b, c, d \rangle : a, b, c \text{ and } p \text{ are binary integers} \\ \text{such that } a^b \equiv c \pmod{p}\}.$$

Show that  $M \in P$ . (Note that the most obvious algorithm does not run in polynomial time. Hint: Try first where  $b$  is a power of 2.)

4. Prove that the following language is undecidable:

$$A = \{\langle M \rangle : M \text{ is a TM with running time } O(n)\}.$$

5. Prove that the following language is undecidable:

$$A = \{\langle M \rangle : L(M) \in \text{TIME}(n)\}.$$