

# Survey of Research Interests for Welcome Week

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# Introduction

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I did have a vague idea of my priorities.

- Something that skirts the boundary between theory-building and applications, being slightly more dominated by theory-building.
- Something that prepares me for a career in academia, but also allows me to entertain job offers in private industry upon graduating.
- Something that has the potential to unify different areas of mathematics. I think of myself of more of a "breadth" mathematician instead of a "depth" mathematician.

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Enter Computational Geometry and Optimization!

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But can we do better?

Using a divide and conquer method, we can achieve  $O(n \log n)$ , and even  $O(n \log \log n)$  deterministic algorithms. If we allow random algorithms, there are algorithms for this problem with expected completion time  $O(n)$ .

# Helly-Type Theorems

When I'm feeling less computational, there are deep theorems that also warrant my attention.

## Theorem (Helly, 1913)

*Let  $X_1, \dots, X_n$  be a finite collection of convex subsets of  $\mathbf{R}^d$ , with  $n > d$ . If the intersection of every  $d + 1$  of these sets is nonempty, then the whole intersection is nonempty, i.e.*

$$\bigcap_{j=1}^n X_j \neq \emptyset$$

Recall that a set is convex if, given any two points in the set, the line connecting them also lies in the set.

# Helly-Type Theorems

This is an amazing result! Even if you gave me  $2^{2^{2^2}}$  convex sets in the plane, I only need to check them in combinations of threes to guarantee non-empty intersection of all of them.

# Example

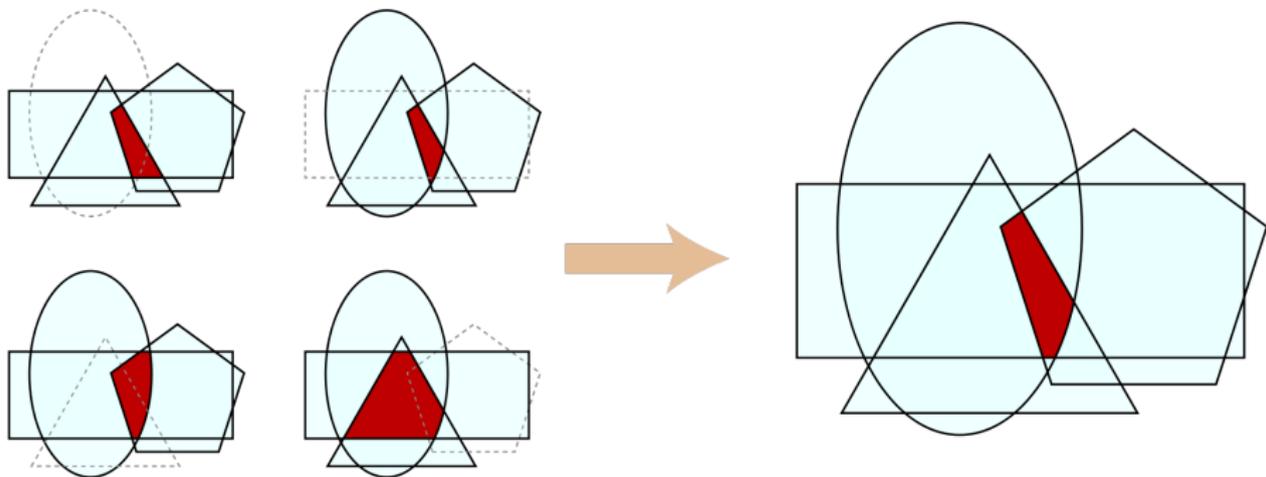


Figure: I shamelessly stole this graphic from wikipedia

# Helly-Type Theorems

This result is so important that it spawned a class of theorems, which are called "Helly-Type".

## Theorem (Helly-Type Theorem)

*If every  $m$  members of a family of objects have property  $P$ , then the entire family has property  $P$*

For the classical Helly's theorem,  $m$  is  $d + 1$ , the objects are convex sets, and  $P$  is the property "have non-empty intersection".

## Doignon's Theorem

Recently, my favorite Helly-Type Theorem has been Doignon's Theorem.

### Theorem (Doignon 1973)

*Consider a polyhedron  $A\mathbf{x} \leq \mathbf{b}$  where  $A$  is a  $d \times n$  matrix and  $\mathbf{b}$  is a vector of  $\mathbf{R}^d$ . If this polyhedron has no integer points in its interior, then there is a subset of rows of  $A$  of cardinality no more than  $2^n$  such that this smaller polyhedron also has no integer points in its interior.*

Note that this is the contrapositive of a Helly-type theorem.

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Clarkson's algorithm is a *Las Vegas Algorithm* which is a type of randomized algorithm that gambles with the speed of completion, but not with the correctness of the solution.

The idea is the following

- Given an integer linear program, there is a subset of the constraints of cardinality  $2^n - 1$  with the same optimal objective value. (Doignon)
- Find it! Use random sampling to quickly grab constraints that are relevant and throw away others.
- Solve this subproblem. It is (hopefully) significantly smaller than the original.

## Recent Developments

From the particular project that I am working on

Theorem (Aliev, De Loera, Louveaux, 2014)

*Given an integer  $k$ , there exists a constant  $c(k, n)$  depending only on the dimension  $n$  and  $k$ , such that if a polyhedron  $\{x : Ax \leq b\}$  contains exactly  $k$  integer solutions, then there exists a subset of the rows of cardinality no more than  $c(k, n)$ , defining a polyhedron that contains exactly the same  $k$  integer solutions.*

This is an amazing result! It extends Doignon's theorem, for which  $k = 0$ , to polyhedra with a fixed number of points inside.

## Open Question

Can Clarkson's Algorithm be extended to return the  $k$  best solutions? The general idea is still the same, but the devil is in the details!

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- This would be extremely helpful in applications.
- Often times it is helpful to consider the " $k$  best numerical options" instead of the just the "best numerical option".
- Implementing this algorithm and testing its success experimentally brings yet another perspective to this project.

Thank you!

Eckhoff, J. (1993), "Helly, Radon, and Carathodory type theorems", Handbook of Convex Geometry A, B, Amsterdam: North-Holland, pp. 389-448.

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M. I. Shamos and D. Hoey. "Closest-point problems." In Proc. 16th Annual IEEE Symposium on Foundations of Computer Science (FOCS), pp. 151-162, 1975