# Computing the Bracket Polynomial of Two and Three Strand Links 

Alexandra Vizcarra<br>Thesis Advisor: Dr. Evgeny Gorskiy



## Contents

1. Introduction ..... 1
2. Background on Knots ..... 2
3. The Bracket Polynomial ..... 3
3.1. Properties ..... 3
3.2. The Jones Polynomial ..... 4
4. Computations of Two Strands ..... 5
5. Temperley-Lieb Matrices ..... 8
6. Computations of Three Strands ..... 9
7. Acknowledgements ..... 11
References ..... 11

## 1. Introduction

We will study invariants of knots and links such as the Kauffman bracket polynomial and the Jones polynomial.

In Sections 2 and 3, we recall the three Reidemeister moves, the definition and the properties of the bracket polynomial and the Jones polynomial following [1]. In particular, the curl relation in Proposition 3.3 describes the behavior of the bracket polynomial under the first Reidemeister move and is used frequently in this paper.
In Section 4, we compute the bracket polynomial for arbitrary $T(2, n)$ torus links which can be obtained as closures of two-strand braids $\sigma_{1}^{n}$. In particular, we get the following result (Theorem 4.2):

$$
B_{n}=\langle T(2, n)\rangle=B_{n}=-A^{n+2}+\frac{-A^{n-6}-(-1)^{n} A^{-2-3 n}}{1+A^{-4}}
$$

In Sections 5 and 6, we take a more abstract approach and define a vector space spanned by all crossingless matchings on 3 strands. This space is called the Temperley-Lieb algebra. Adding a crossing on top of a crossingless matching leads to a matrix which we compute in Section 6. There are two such matrices $T_{1}, T_{2}$ corresponding to the generators of the braid group. We verify that the braid relation $T_{1} T_{2} T_{1}=T_{2} T_{1} T_{2}$ holds for such matrices.
Finally, in Theorem 6.6, we compute a matrix for the $(3,3 k)$ torus link by computing the powers $\left(T_{1} T_{2}\right)^{3 k}$ inductively. This allows us to prove Theorem 6.7. which gives an explicit formula for the bracket polynomial of $(3,3 k)$ torus link for all $k \geq 0$.

## 2. Background on Knots

Definition 2.1. A Reidemeister move is one of three ways to change the projection of a knot that will change the relation between the crossings.

Definition 2.2. The first Reidemeister move allows one to put in or take out a twist in the knot.


Definition 2.3. The second Reidemeister move allows one to either add two crossings or remove two crossings.


Definition 2.4. The third Reidemeister move allows one to slide a strand of the knot from one side of a crossing to the other side of the crossing.


Theorem 2.5. Two diagrams correspond to the same knot or link if and only if they are related by a sequence of Reidemeister moves.

For more details see [1].
We will also need some facts about the braid group. The braid group on $n$ strands has generators $\sigma_{1}, \ldots, \sigma_{n-1}$ and relations

$$
\begin{equation*}
\sigma_{i} \sigma_{i+1} \sigma_{i}=\sigma_{i+1} \sigma_{i} \sigma_{i+1}, \sigma_{i} \sigma_{j}=\sigma_{j} \sigma_{i}(|i-j| \geq 2) \tag{1}
\end{equation*}
$$

The second Reidemeister move corresponds to the relation $\sigma_{i} \sigma_{i}^{-1}=1$ and the third Reidemeister move corresponds to the first equation in (1).

## 3. The Bracket Polynomial

Definition 3.1. The bracket polynomial is defined by the following relations:

$$
\searrow / \=)\left(\cdot \mathrm{A}+\widetilde{\sim} \cdot \mathrm{A}^{-1}\right.
$$

and


### 3.1. Properties.

Proposition 3.2. The Bracket polynomial does not change under the second and third Reidemeister moves.

See Section 6.1 in [1].

Proposition 3.3. The Bracket polynomial does change under the first Reidemeister move, so a Curl Revalion is used. The bracket polynomials before and after adding a curl differ by a factor $-A^{ \pm 3}$.

Proof. The definition of the first Reidemeister move is given by


Apply 3.1

$$
\begin{aligned}
\mathrm{O} & \rightarrow O \cdot A+\stackrel{\backsim}{\sim} A^{-1} \\
& \left.=) \cdot\left(-A^{2}-A^{-2}\right) \cdot A+\right) \cdot A^{-1} \\
& =) \cdot\left(-A^{3}-A^{-1}+A^{-1}\right) \\
& =) \cdot-A^{3}
\end{aligned}
$$

$$
\begin{aligned}
& =) \cdot\left(A+\left(-A^{2}-A^{-2}\right) A^{-1}\right) \\
& =) \cdot-\mathrm{A}^{-3}
\end{aligned}
$$

Thus, we obtain the Curl Relation $-A^{ \pm 3}$.

### 3.2. The Jones Polynomial.

How do we fix the issue with curls, and define an actual link invariant which does not change under all Reidemeister moves? We orient the link by choosing a direction for each component.

Definition 3.4. A positive crossing ( +1 ) is defined by


Definition 3.5. A negative crossing (-1) is defined by


Definition 3.6. The writhe of an oriented link $w(L)$ is the difference between the number of positive and negative crossings, or

$$
w\langle L\rangle=\text { (number of positive crossings) }-(\text { number of negative crossings })
$$

Example 3.7. Orientation of the Hopf Link:


Since there are two negative crossings, we have $w\langle L\rangle=-2$.

Example 3.8. Orientation of the Trefoil Knot:


Since there are three positive crossings, we have $w\langle L\rangle=3$.

Definition 3.9. Finally, we can define the Jones polynomial as

$$
J(L)=\left(-A^{3}\right)^{-w(L)}\langle L\rangle
$$

Theorem 3.10. $J(L)$ is a link invariant and does not change under all three Reidemeister moves

## 4. Computations of Two Strands

In this section we consider the $T(2, n)$ torus links which can be obtained as closures of two-strand braids $\sigma_{1}^{n}$. Let $B_{n}$ be the bracket polynomial for $T(2, n)$.

Lemma 4.1. The polynomials $B_{n}$ satisfy the recursion

$$
B_{n}=A B_{n-1}+A^{-1}\left(-A^{-3}\right)^{n-1}
$$

Proof. We present the following link $B_{n}$ with $n$ crossings:


We first apply the equation in Definition 3.1 to the bottom crossing, which gives us


Notice that
 is the original knot, but with $n-1$ crossings. We call this $B_{n-1}$.

. We will eventually obtain the unknot if we
We can also note that we can apply the curl relation to continuously apply the curl relation $(n-1)$ times to this knot.

Thus, we have obtained the recursive formula $B_{n}=A B_{n-1}+A^{-1}\left(-A^{-3}\right)^{n-1}$

Theorem 4.2. The bracket polynomial for the $T(2, n)$ torus link is

$$
B_{n}=-A^{n+2}+\frac{-A^{n-6}-(-1)^{n} A^{-2-3 n}}{1+A^{-4}}
$$

Proof. We prove this by induction, using the previous lemma.
For the base case : $n=0$, we have an unlink with two components, so $B_{0}=-A^{2}-A^{-2}$.
On the other hand, we have

$$
-A^{2}+\frac{-A^{-6}-(-1)^{0} A^{-2}}{1+A^{-4}}=-A^{2}-\frac{A^{-2}\left(1+A^{-4}\right)}{1+A^{-4}}=-A^{2}-A^{-2}
$$

For the inductive step, we assume that the $n-1$ case is true:

$$
B_{n-1}=-A^{n+1}+\frac{-A^{n-7}-(-1)^{n-1} \cdot A^{-2-3(n-1)}}{1+A^{-4}}
$$

So,

$$
\begin{aligned}
A B_{n-1}+A^{-1} \cdot\left(-A^{-3}\right)^{n-1} & =A B_{n-1}+A^{-1} \cdot(-1)^{n-1} \cdot A^{-3 n+3} \\
& =A B_{n-1}+(-1)^{n-1} \cdot A^{-3 n+2}
\end{aligned}
$$

Then we plug in $B_{n-1}$, which gives us

$$
-A^{n+2}+\frac{-A^{n-6}-(-1)^{n-1} \cdot A^{-2-3(n-1)+1}}{1+A^{-4}}+(-1)^{n-1} \cdot A^{-3 n+2}
$$

Simplifying this expression would give us

$$
\begin{gathered}
-A^{n+2}+\frac{-A^{n-6}-(-1)^{n-1} \cdot A^{-3 n+2}+(-1)^{n-1} \cdot A^{-3 n+2}+(-1)^{n-1} \cdot A^{-3 n+2-4}}{1+A^{-4}} \\
=-A^{n+2}+\frac{-A^{n-6}+(-1)^{n-1} \cdot A^{-2-3 n}}{1+A^{-4}} \\
=-A^{n+2}+\frac{-A^{n-6}-(-1)^{n} \cdot A^{-2-3 n}}{1+A^{-4}}
\end{gathered}
$$

Thus, the induction proof is completed.

Example 4.3. For $n=1$ we get the unknot:


When we apply the curl relation, we get $B_{1}=-A^{3}$. We can also plug in $n=1$ into Theorem 4.2 and get

$$
B_{1}=-A^{3}+\frac{-A^{-5}+A^{-5}}{1+A^{-4}}=-A^{3}
$$

Example 4.4. For $n=2$ we get the Hopf link


We apply Definition 3.1 to the bottom crossing to get
?

First, we look at $(3.1$ and apply to get

$$
\left.(\because)=(\Omega) \cdot \mathrm{Q}+(\square) \cdot \mathrm{A}^{-1}\right) \cdot \mathrm{A}
$$

The equation we end up with is $\left(\left(-A^{2}-A^{-2}\right) \cdot A+1 \cdot A^{-1}\right) \cdot A$, which then simplifies to $-A^{4}$

Then we apply the curl relation to
 and we get

$$
\square=-A^{-3} \cdot A^{-1}
$$

which simplifies to $-A^{-4}$. Thus, the polynomial equals

$$
B_{2}=-A^{4}-A^{-4}
$$

We can also plug $n=2$ into Theorem 4.2 and get

$$
B_{2}=-A^{4}+\frac{-A^{-4}-A^{-8}}{1+A^{-4}}=-A^{4}-A^{-4}
$$

Example 4.5. For $n=3$ we get the trefoil knot


We can follow the same process as the previous example where we apply 3.1 to the bottom crossing. Even better, we can use 4.1 to obtain the polynomial
Since there are three crossings, we have

$$
B_{3}=A B_{2}+A^{-1}\left(-A^{-3}\right)^{2}
$$

We found that $B_{2}=-A^{4}-A^{-4}$ in the previous example. Now,

$$
B_{3}=A\left(-A^{4}-A^{-4}\right)+A^{-1}\left(-A^{-3}\right)^{2}
$$

Thus, the polynomial equals $B_{3}=-A^{5}-A^{3}+A^{-7}$.
We can also plug $n=3$ into Theorem 4.2 and get

$$
\begin{aligned}
B_{3} & =-A^{5}+\frac{-A^{-3}+A^{-11}}{1+A^{-4}} \\
& =-A^{5}-A^{-3} \frac{1-A^{-8}}{1+A^{-4}} \\
& =-A^{5}-A^{-3}\left(1-A^{-4}\right) \\
& =-A^{5}-A^{3}+A^{-7}
\end{aligned}
$$

5. Temperley-Lieb Matrices

Let $V$ be the 5 dimensional space spanned by the pictures in Figure 1 also known as crossingless matchings.

$$
\mathrm{e}_{1}=\left|\left|\left|\mathrm{e}_{2}=\circlearrowright\right| \quad \mathrm{e}_{3}=\right|\right.
$$



Figure 1.

Define $T_{1}$ to be the linear operator which adds the crossing $\sigma_{1}$ on top of a diagram in $V$ :
and define $T_{2}$ be the linear operator which adds the crossing $\sigma_{2}$ on top:
We can use the defining relation for the bracket polynomial to write the matrix for $T_{1}$ in the basis of $V$, see Figure 2

$$
=
$$

Figure 2.

## 6. Computations of Three Strands

Definition 6.1. We define matrices $T_{1}$ and $T_{2}$ by the following equations

$$
\begin{aligned}
T_{1} & =\left(\begin{array}{ccccc}
A & 0 & 0 & 0 & 0 \\
A^{-1} & -A^{-3} & 0 & 0 & A^{-1} \\
0 & 0 & A & 0 & 0 \\
0 & 0 & A^{-1} & -A^{-3} & 0 \\
0 & 0 & 0 & 0 & A
\end{array}\right) \\
T_{2} & =\left(\begin{array}{ccccc}
A & 0 & 0 & 0 & 0 \\
0 & A & 0 & 0 & 0 \\
A^{-1} & 0 & -A^{-3} & A^{-1} & 0 \\
0 & 0 & 0 & A & 0 \\
0 & A^{-1} & 0 & 0 & -A^{-3}
\end{array}\right)
\end{aligned}
$$

It is easy to check the following formula for the product $T_{1} T_{2}$.
Proposition 6.2. The product of $T_{1} T_{2}$ is the matrix

$$
T_{1} T_{2}=\left(\begin{array}{ccccc}
A^{2} & 0 & 0 & 0 & 0 \\
1 & 0 & 0 & 0 & -A^{-4} \\
1 & 0 & -A^{-2} & 1 & 0 \\
A^{-2} & 0 & -A^{-4} & 0 & 0 \\
0 & 1 & 0 & 0 & -A^{-2}
\end{array}\right)
$$

To every braid on 3 strands, we associate a matrix by replacing $\sigma_{1}$ and $\sigma_{2}$ with $T_{1}$ and $T_{2}$.
The following proposition shows that $T_{1}$ and $T_{2}$ satisfy braid relations (1), and therefore this assignment of matrices to braids is well defined.

Proposition 6.3. $T_{1} T_{2} T_{1}=T_{2} T_{1} T_{2}$

Proof. We compute this by matrix multiplication.

$$
\begin{aligned}
& T_{1} T_{2} T_{1}=\left(\begin{array}{ccccc}
A^{3} & 0 & 0 & 0 & 0 \\
A & 0 & 0 & 0 & -A^{-3} \\
A & 0 & 0 & -A^{-3} & 0 \\
A^{-1} & 0 & -A^{-3} & 0 & 0 \\
A^{-1} & -A^{-3} & 0 & 0 & 0
\end{array}\right) \\
& T_{2} T_{1} T_{2}=\left(\begin{array}{ccccc}
A^{3} & 0 & 0 & 0 & 0 \\
A & 0 & 0 & 0 & -A^{-3} \\
A & 0 & 0 & -A^{-3} & 0 \\
A^{-1} & 0 & -A^{-3} & 0 & 0 \\
A^{-1} & -A^{-3} & 0 & 0 & 0
\end{array}\right)
\end{aligned}
$$

Thus, $T_{1} T_{2} T_{1}=T_{2} T_{1} T_{2}$

Example 6.4. The figure eight knot is obtained as the closure of the braid $\sigma_{1} \sigma_{2}^{-1} \sigma_{1} \sigma_{2}^{-1}$ and corresponds to the matrix

$$
T_{1} T_{2}^{-1} T_{1} T_{2}^{-1}=\left(\begin{array}{ccccc}
1 & 0 & 0 & 0 & 0 \\
\frac{2 A^{4}-1}{A^{6}} & \frac{-A^{12}+\left(A^{4}-1\right)^{2}}{A^{8}} & 0 & 0 & \frac{A^{8}-A^{4}+1}{A^{2}} \\
2 A^{2}-A^{6} & 0 & A^{8}-A^{4} & \frac{-A^{8}+A^{4}-1}{A^{2}} & 0 \\
\frac{-A^{8}+2 A^{4}-1}{A^{4}} & 0 & \frac{A^{8}-A^{4}+1}{A^{2}} & \frac{-A^{12}+\left(A^{4}-1\right)^{2}}{A^{8}} & 0 \\
1 & \frac{-A^{8}+A^{4}-1}{A^{2}} & 0 & 0 & A^{8}-A^{4}
\end{array}\right)
$$

The $(3, m)$ torus links are obtained as the closures of the braids $\left(\sigma_{1} \sigma_{2}\right)^{m}$. The corresponding matrices are $\left(T_{1} T_{2}\right)^{m}$.

Example 6.5. The braid $\left(\sigma_{1} \sigma_{2}\right)^{3}$ is a full twist, or a $360^{\circ}$ rotation. The corresponding matrix is

$$
\left(T_{1} T_{2}\right)^{3}=\left(\begin{array}{ccccc}
A^{6} & 0 & 0 & 0 & 0 \\
\frac{A^{8}-1}{A^{4}} & A^{-6} & 0 & 0 & 0 \\
\frac{A^{8}-1}{A^{4}} & 0 & A^{-6} & 0 & 0 \\
\frac{A^{A^{4}-1}}{A^{2}} & 0 & 0 & A^{-6} & 0 \\
\frac{A^{4}-1}{A^{2}} & 0 & 0 & 0 & A^{-6}
\end{array}\right)
$$

Theorem 6.6. For $m=3 k$ we have

$$
\left(T_{1} T_{2}\right)^{3 k}=\left(\begin{array}{ccccc}
A^{6 k} & 0 & 0 & 0 & 0  \tag{2}\\
x_{k} & A^{-6 k} & 0 & 0 & 0 \\
x_{k} & 0 & A^{-6 k} & 0 & 0 \\
y_{k} & 0 & 0 & A^{-6 k} & 0 \\
y_{k} & 0 & 0 & 0 & A^{-6 k}
\end{array}\right)
$$

where
and

$$
x_{k}=\frac{\left(1+A^{12}+A^{24}+A^{36}+\cdots+A^{12(k-1)}\right)\left(A^{8}-1\right)}{A^{6 k-2}}
$$

$$
y_{k}=\frac{\left(1+A^{12}+A^{24}+A^{36}+\cdots+A^{12(k-1)}\right)\left(A^{4}-1\right)}{A^{6 k-4}}
$$

Proof. We prove by induction in $k$. The base case $k=1$ is covered by Example 6.5. To prove the inductive step, assume that $\left(T_{1} T_{2}\right)^{3 k}$ satisfies (2). Then

$$
\left(T_{1} T_{2}\right)^{3(k+1)}=\left(T_{1} T_{2}\right)^{3 k}\left(T_{1} T_{2}\right)^{3}
$$

and by multiplying the matrices, we get the recursion relations

$$
x_{k+1}=A^{6} x_{k}+A^{-6 k} \frac{A^{8}-1}{A^{4}}, y_{k+1}=A^{6} y_{k}+A^{-6 k} \frac{A^{4}-1}{A^{2}}
$$

It remains to prove that $x_{k}$ and $y_{k}$ are given by the above formulas.
Now, let's prove the formula for $x_{k}$.
Base case $k=1$ :

$$
x_{1}=\frac{\left(A^{8}-1\right)}{A^{4}}
$$

For the inductive step, assume that the formula for $x_{k}$ is true.

$$
\begin{aligned}
x_{k+1} & =\frac{\left(1+A^{12}+\cdots+A^{12(k-1)}\right)\left(A^{8}-1\right)}{A^{6 k-2}} \cdot A^{6}+\frac{\left(A^{8}-1\right)}{A^{4}} \cdot A^{-6 k} \\
& =\left(A^{8}-1\right)\left(\frac{\left(1+A^{12}+\cdot+A^{12(k-1)}\right) \cdot A^{6} \cdot A^{6}}{A^{6 k-2} \cdot A^{6}}+\frac{1}{A^{6 k-2} \cdot A^{6}}\right) \\
& =\left(A^{8}-1\right)\left(\frac{A^{12}+A^{24}+\cdot+A^{12 k}+1}{A^{6 k-2} \cdot A^{6}}\right)
\end{aligned}
$$

The formula for $y_{k}$ is proven similarly.
Thus, the proof is completed.

Theorem 6.7. The bracket polynomial for the torus link $T(3,3 k)$ is given by the equation

$$
\langle T(3,3 k)\rangle=A^{6 k}\left(-A^{2}-A^{-2}\right)^{2}+2 x_{k}\left(-A^{2}-A^{-2}\right)+2 y_{k}
$$

where $x_{k}$ and $y_{k}$ are as in Theorem 6.6.
Proof. First, we express $\left(T_{1} T_{2}\right) e_{1}$ in the basis of $V$ using Theorem 6.6 .

$$
\begin{equation*}
\left(T_{1} T_{2}\right)^{3 k} \cdot e_{1}=A^{6 k} e_{1}+x_{k} e_{2}+x_{k} e_{3}+y_{k} e_{4}+y_{k} e_{5} \tag{3}
\end{equation*}
$$

Then, we compute the closure of each $e_{i}$ for $e=1, \cdots, 5$ by using Rules 1 and 2






Now the bracket polynomial of $T(3,3 k)$ can be obtained as follows:

$$
\langle T(3,3 k)\rangle=\left\langle\left(T_{1} T_{2}\right)^{3 k} \cdot e_{1}\right\rangle=A^{6 k}\left\langle e_{1}\right\rangle+x_{k}\left\langle e_{2}\right\rangle+x_{k}\left\langle e_{3}\right\rangle+y_{k}\left\langle e_{4}\right\rangle+y_{k}\left\langle e_{5}\right\rangle
$$

We then use the values for $\left\langle e_{i}\right\rangle$ to obtain the formula given in Theorem.

## 7. Acknowledgements

I would like to thank my thesis advisor, Dr. Evgeny Gorskiy, for introducing this topic to me, guiding me through research and writing, and for being such a wonderful professor. I would also like to thank my peers and TAs for providing me with LaTex help and for proof-reading this thesis.

## References

[1] C. Adams. The Knot Book.

